

# Hash-based Signatures

IETF/IRTF CFRG Draft on XMSS



Fraunhofer Workshop Series 01 – Post-Quantum Cryptography in Practice  
Speaker: Dr. Bernhard Jungk

# eXtended Merkle Signature Scheme

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Why should we look into XMSS?

Hash-based signatures have many advantages:

- Based on well understood security notions
  - » Cryptographic hash functions are hard to invert, **also for quantum computers**
  - » Merkle trees well studied since the 1980ies
- Hash functions are well understood (especially after SHA-3 competition)
- Fast signing and verification operations possible
- Relatively easy to understand and implement

# eXtended Merkle Signature Scheme

Why should we look into XMSS?

XMSS is a promising candidate for

- Applications with relatively low amount of signatures
- One- or many-times firmware updates
- Digital signatures for documents (e.g. contracts, email)
- Long-term archival of important digital assets
- PKI Certificates (e.g. Root CA)

# eXtended Merkle Signature Scheme

Why should we look into XMSS?

IRTF is part of IETF

- Oriented towards research and long-term trends

Important trend – PQC

- Quantum computer attacks are likely
- Design of replacements for traditional public key crypto

Standardization needed

- Interoperability
- Implementation Guidelines

# eXtended Merkle Signature Scheme

Our Contribution

## Implementation experience

- Benchmarking against other schemes
- Learn good trade-offs for different application scenarios, cost reductions, side-channels, etc.

Target Platform: Hardware, i.e. FPGAs and ASICs

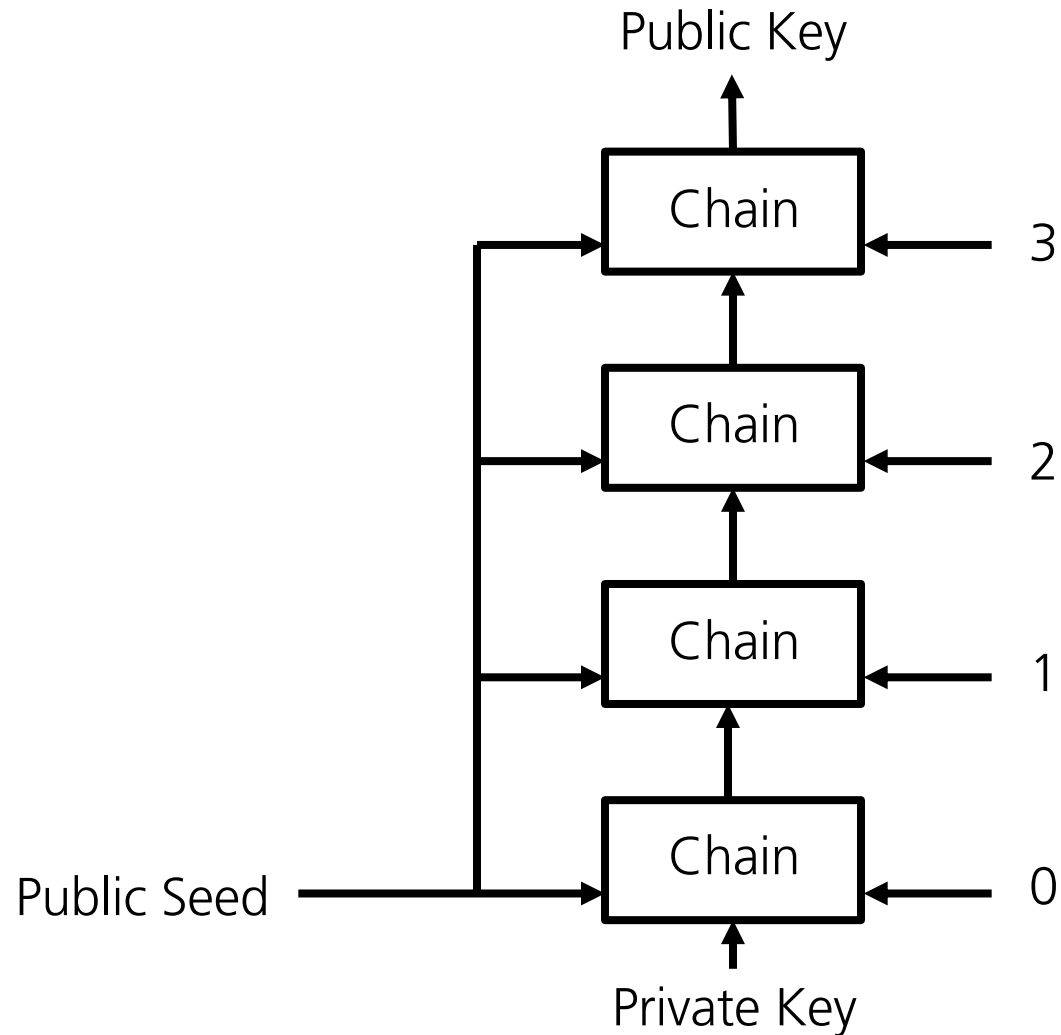
Cooperation:

- Yale University in New Haven, US
- Fraunhofer SIT in Darmstadt, Germany
- Fraunhofer Singapore

# Recap Winternitz One-Time Signatures

# Winternitz One-Time Scheme+

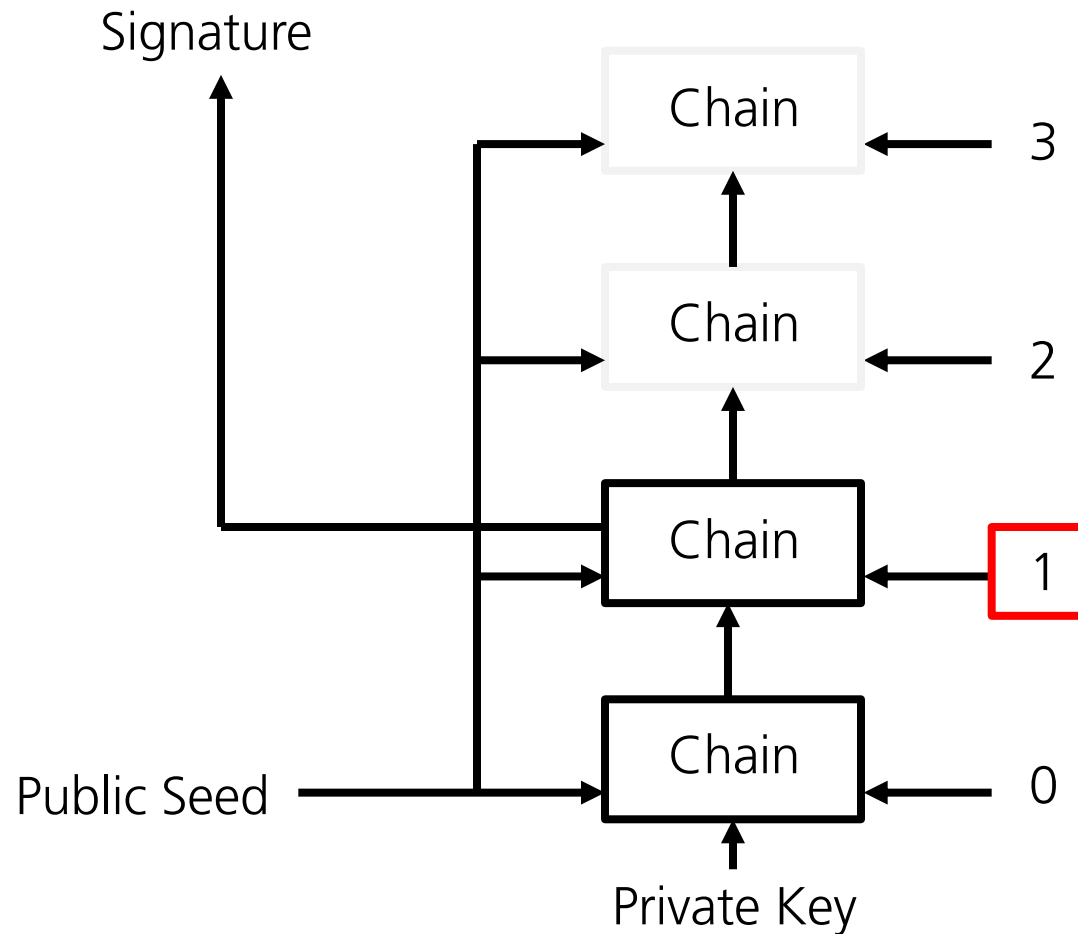
Basic Principle – Public Key Generation





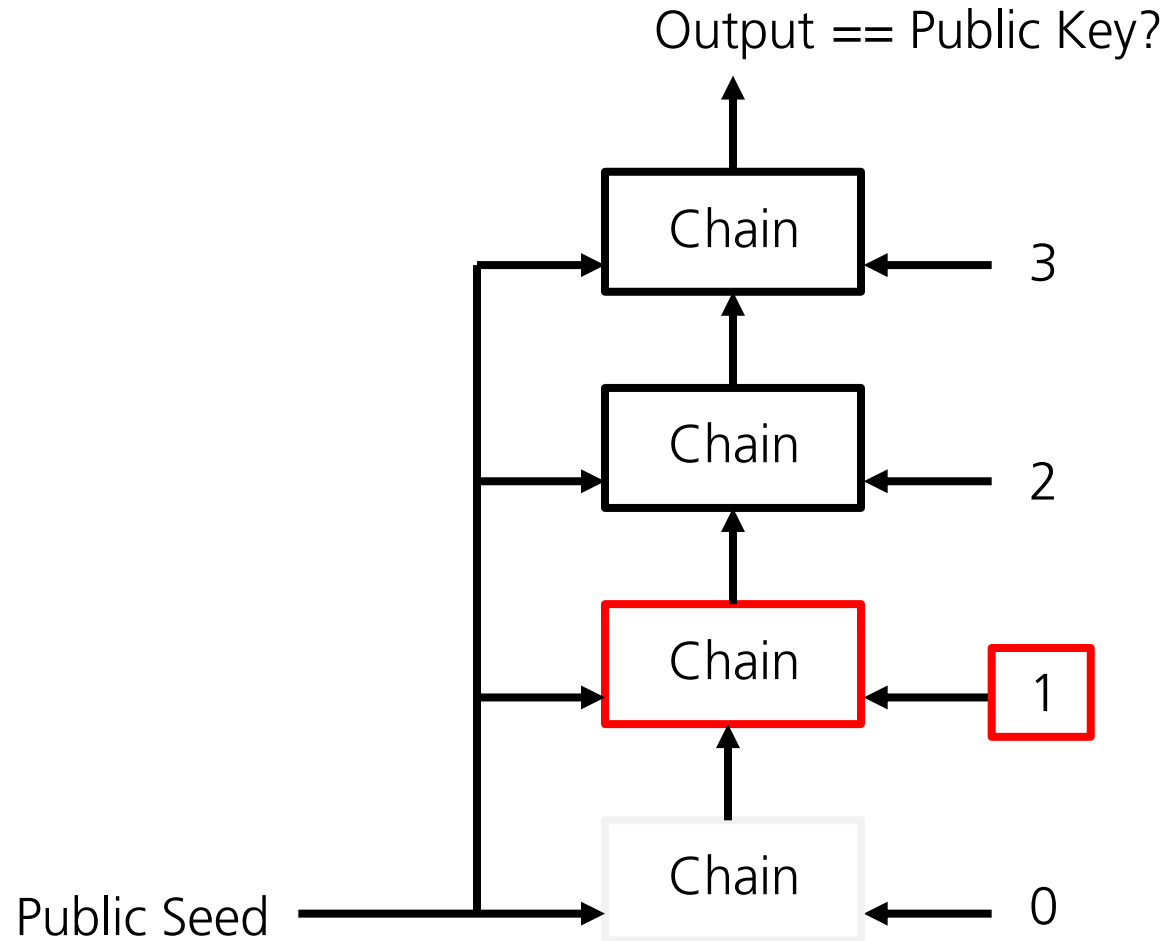
# Winternitz One-Time Scheme+

Basic Principle – Signature Generation



# Winternitz One-Time Scheme+

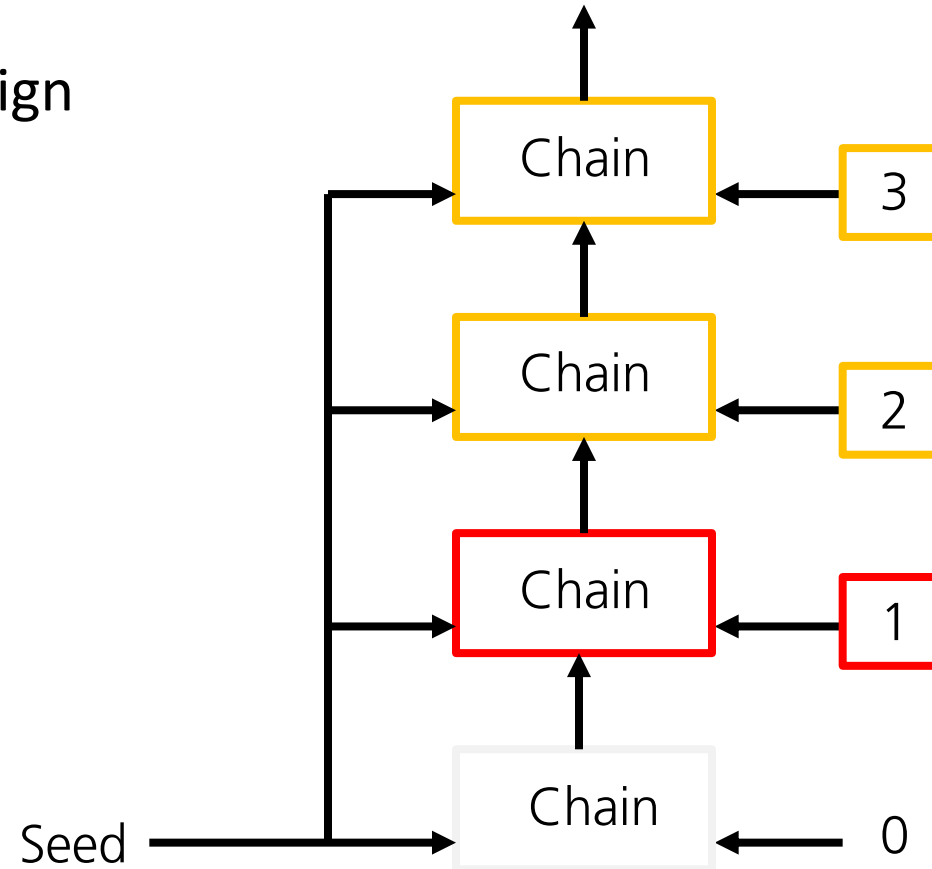
Basic Principle – Signature Verification



# Winternitz One-Time Scheme+

## Basic Principle

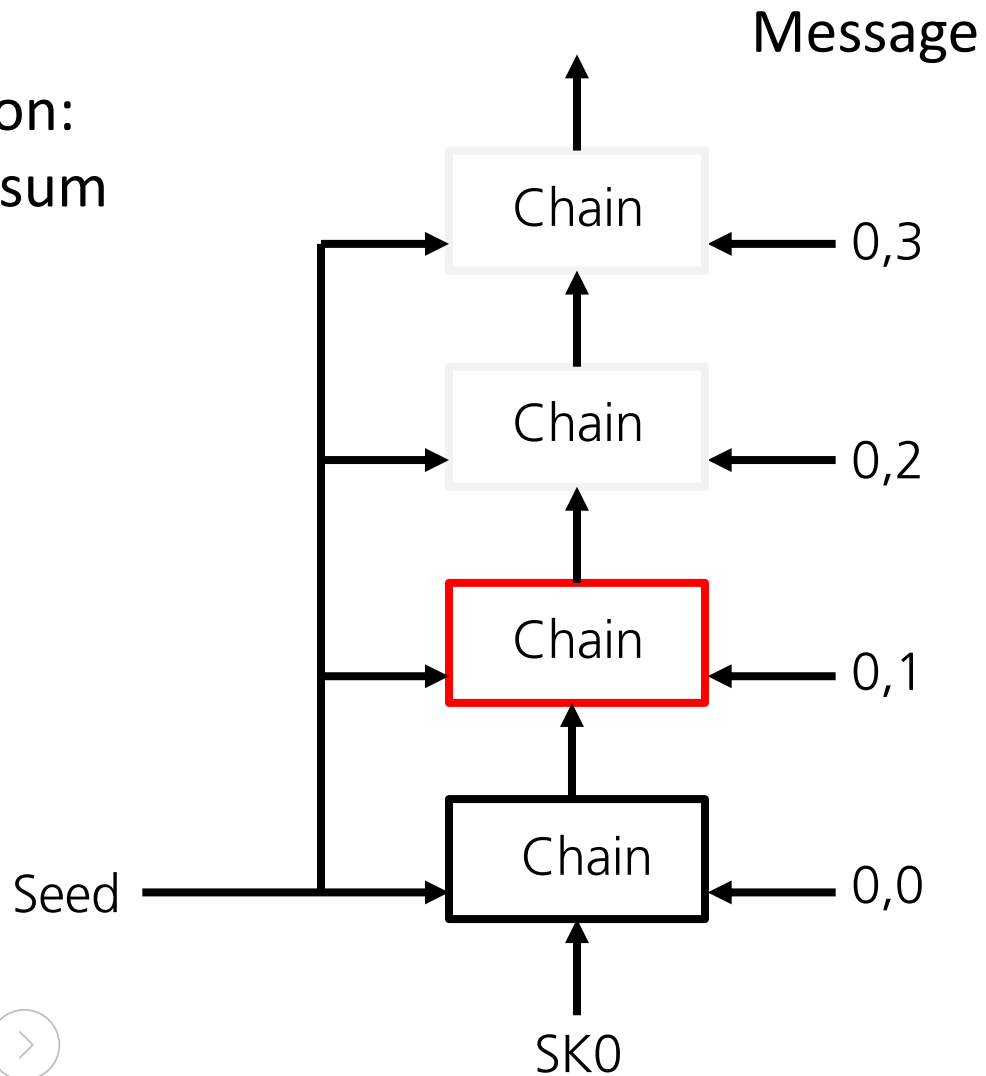
Problem:  
Signer reveals how to sign  
other messages with  
the same key



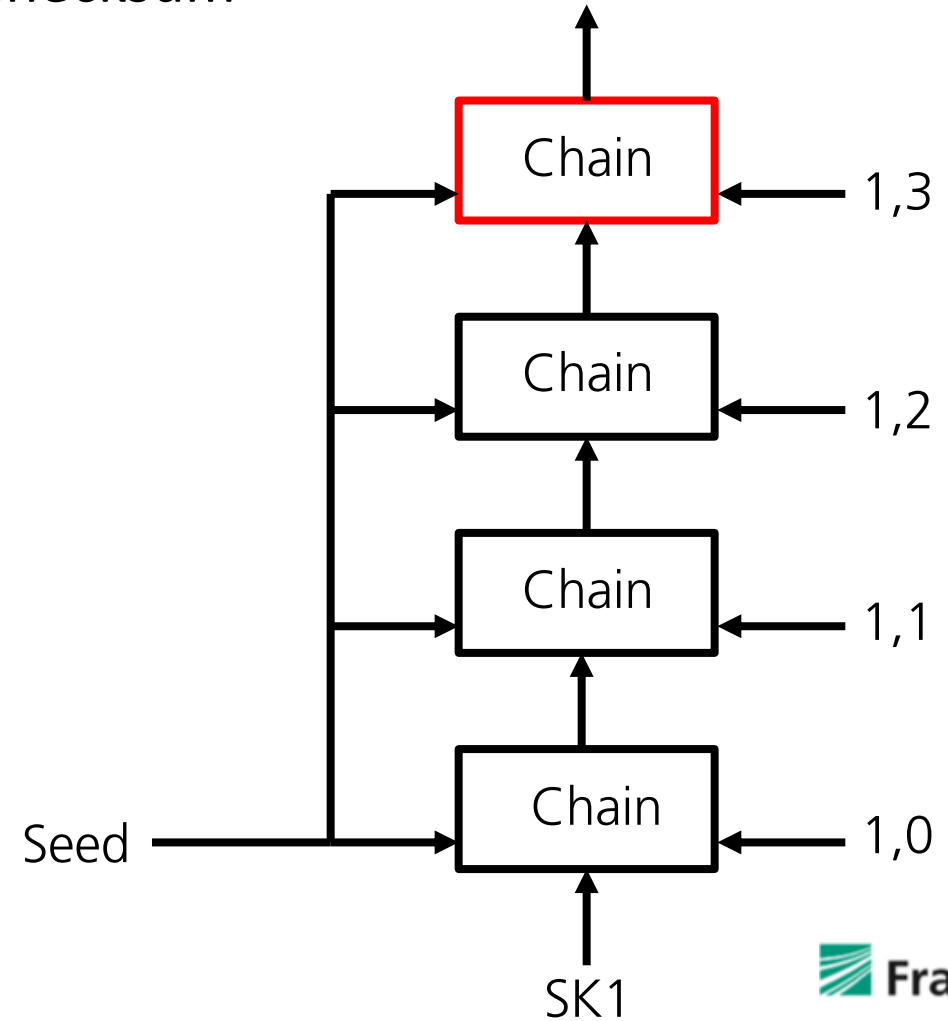
# Winternitz One-Time Scheme+

Basic Principle

Solution:  
Checksum

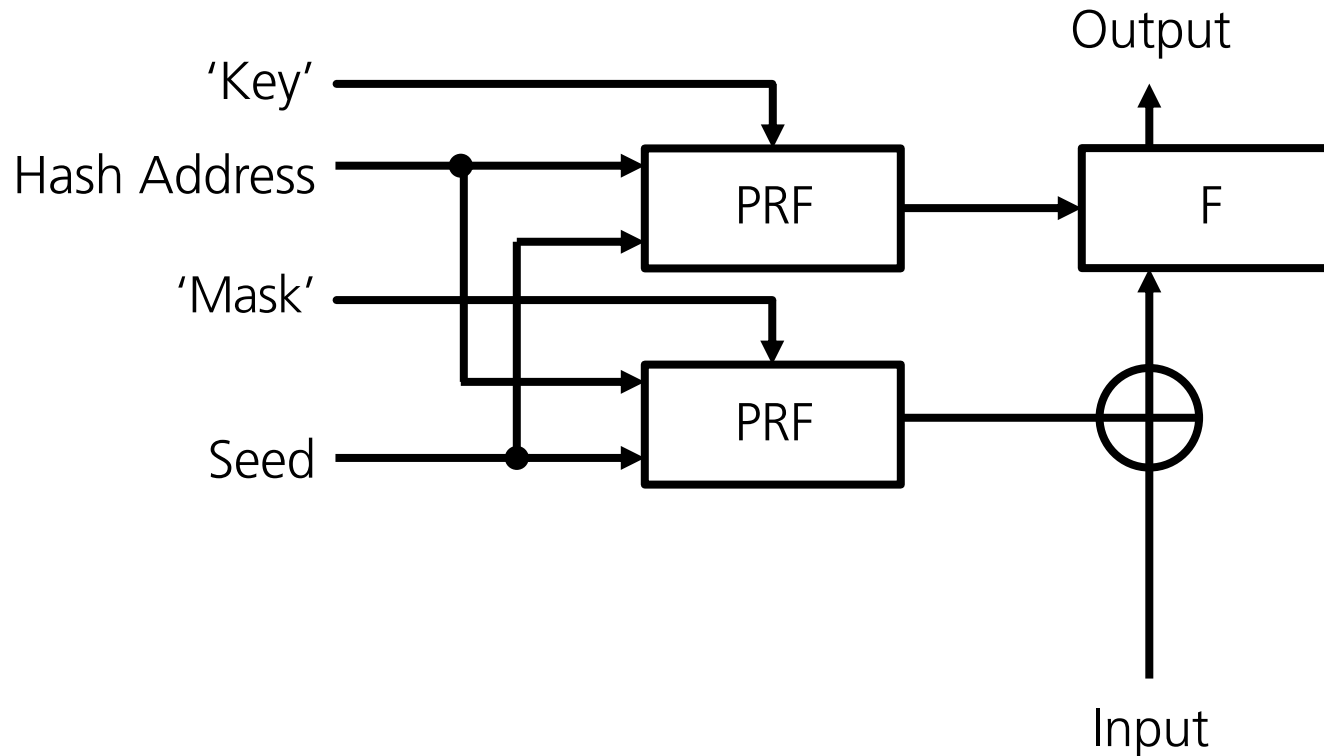


Checksum



# Winternitz One-Time Scheme+

Chaining Function for XMSS

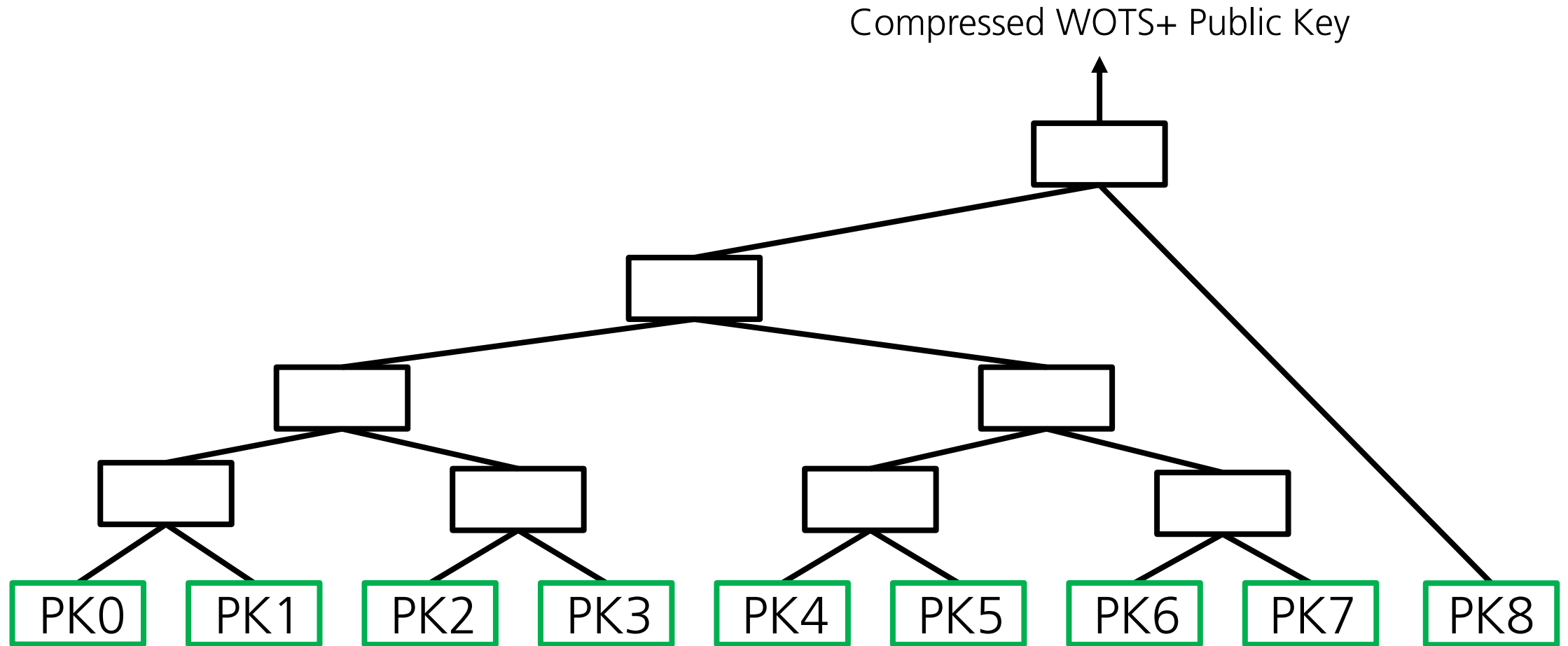


PRF – Pseudorandom function  
F – Keyed hash function

# eXtended Merkle Signature Scheme

# eXtended Merkle Signature Scheme

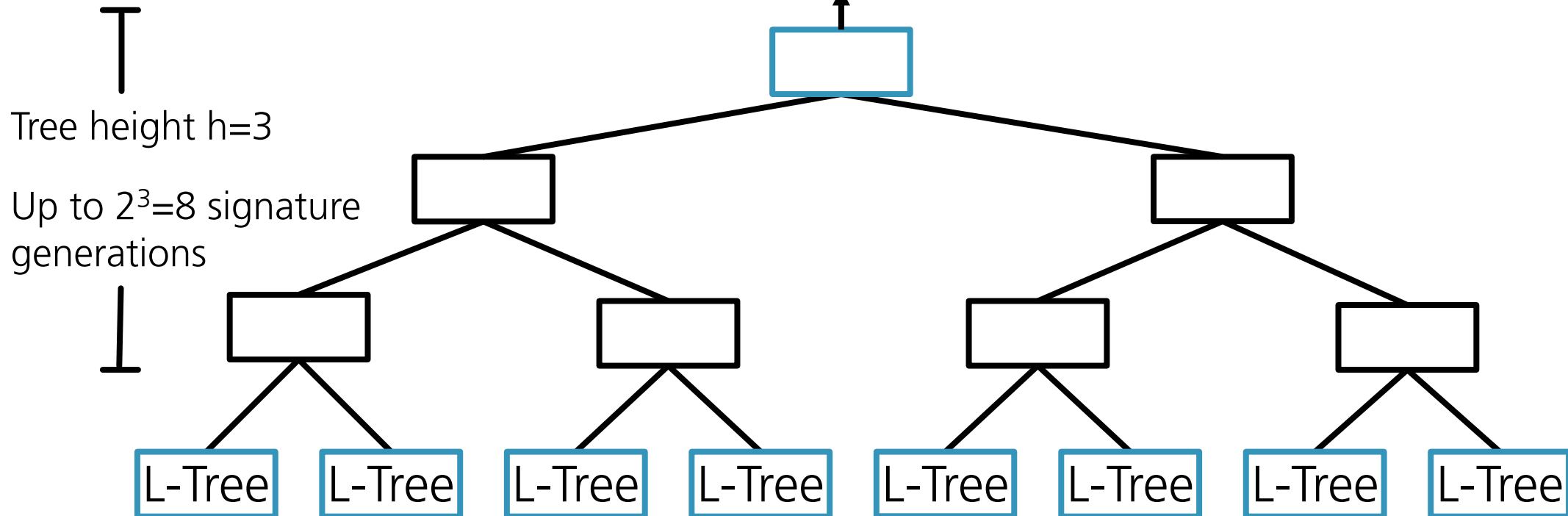
L-Tree – Public Key Generation



# eXtended Merkle Signature Scheme

XMSS Tree – Public Key Generation

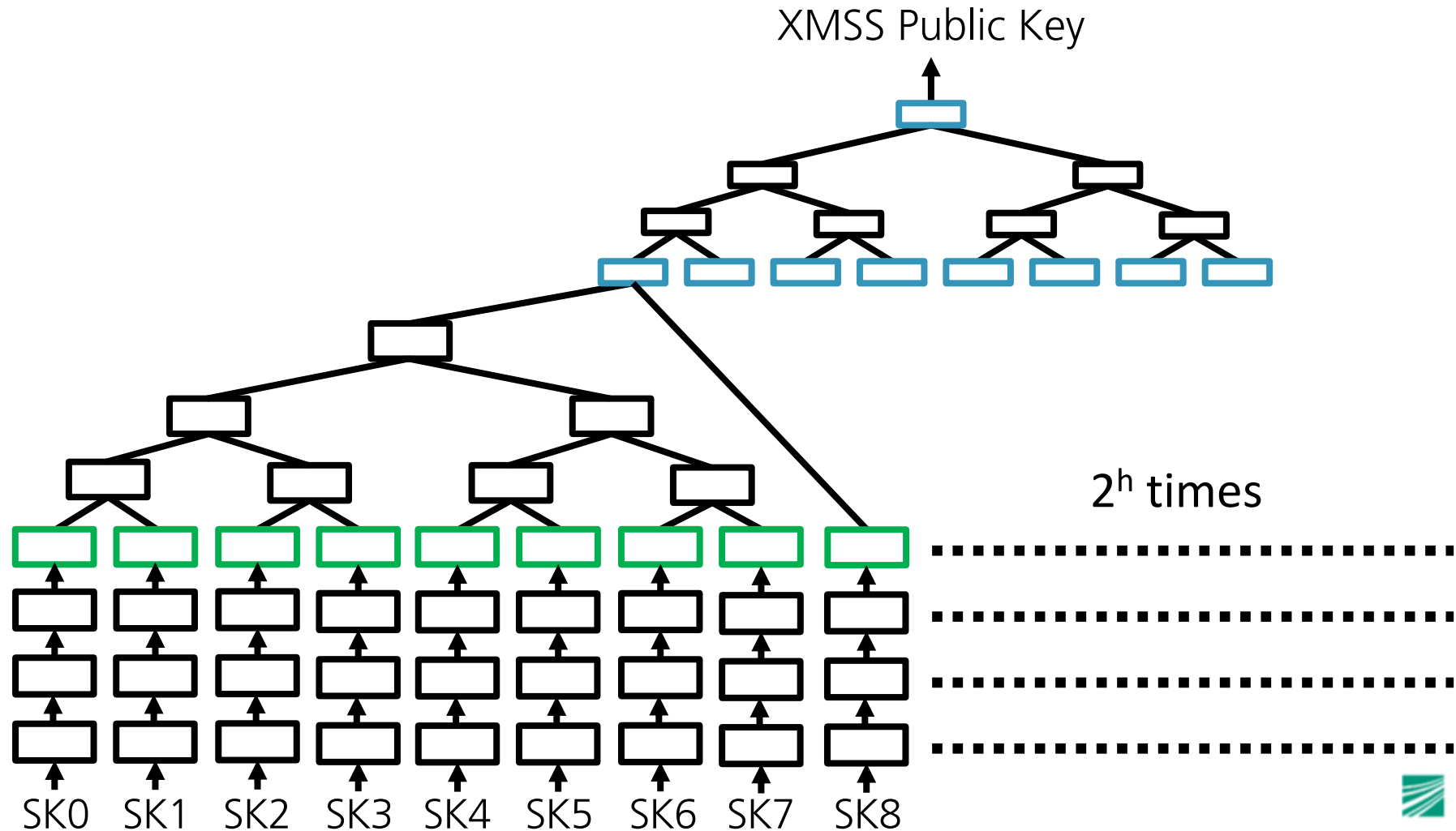
XMSS Public Key





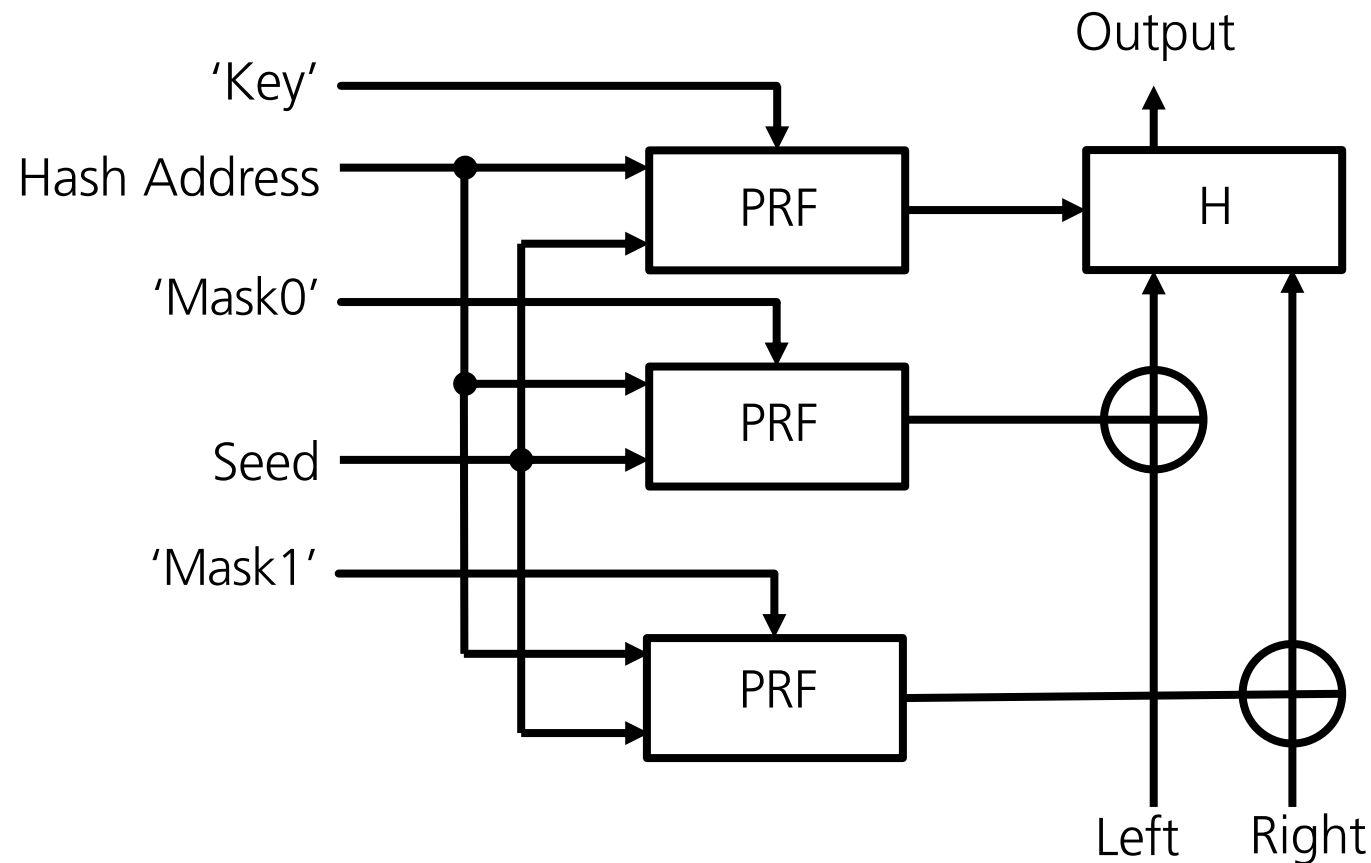
# eXtended Merkle Signature Scheme

The Complete Picture – Public Key Generation



# eXtended Merkle Signature Scheme

rand\_hash



PRF – Pseudorandom function  
H – Keyed hash function

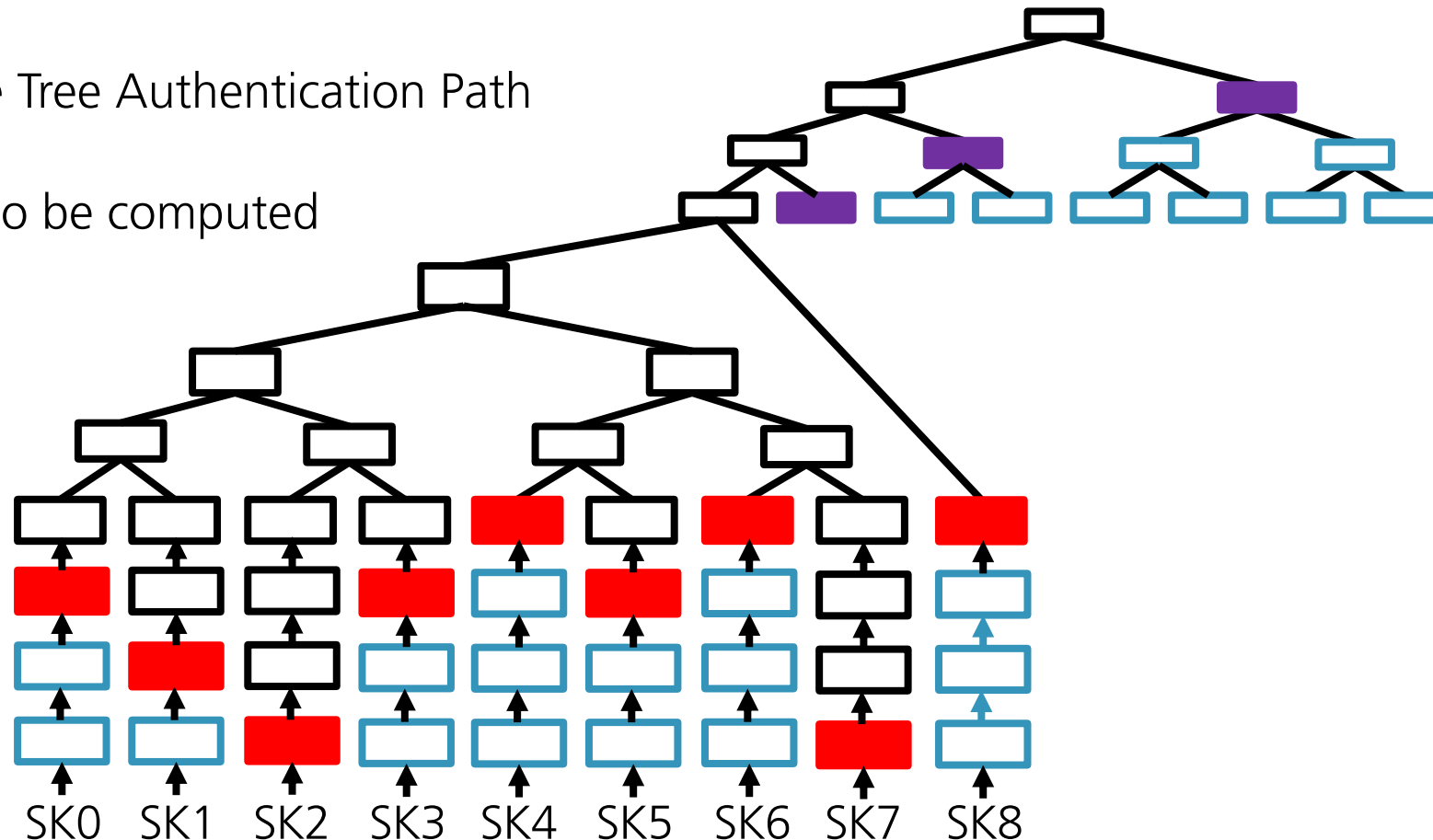
# eXtended Merkle Signature Scheme

Signature Generation – Message 1

 WOTS+ Signature

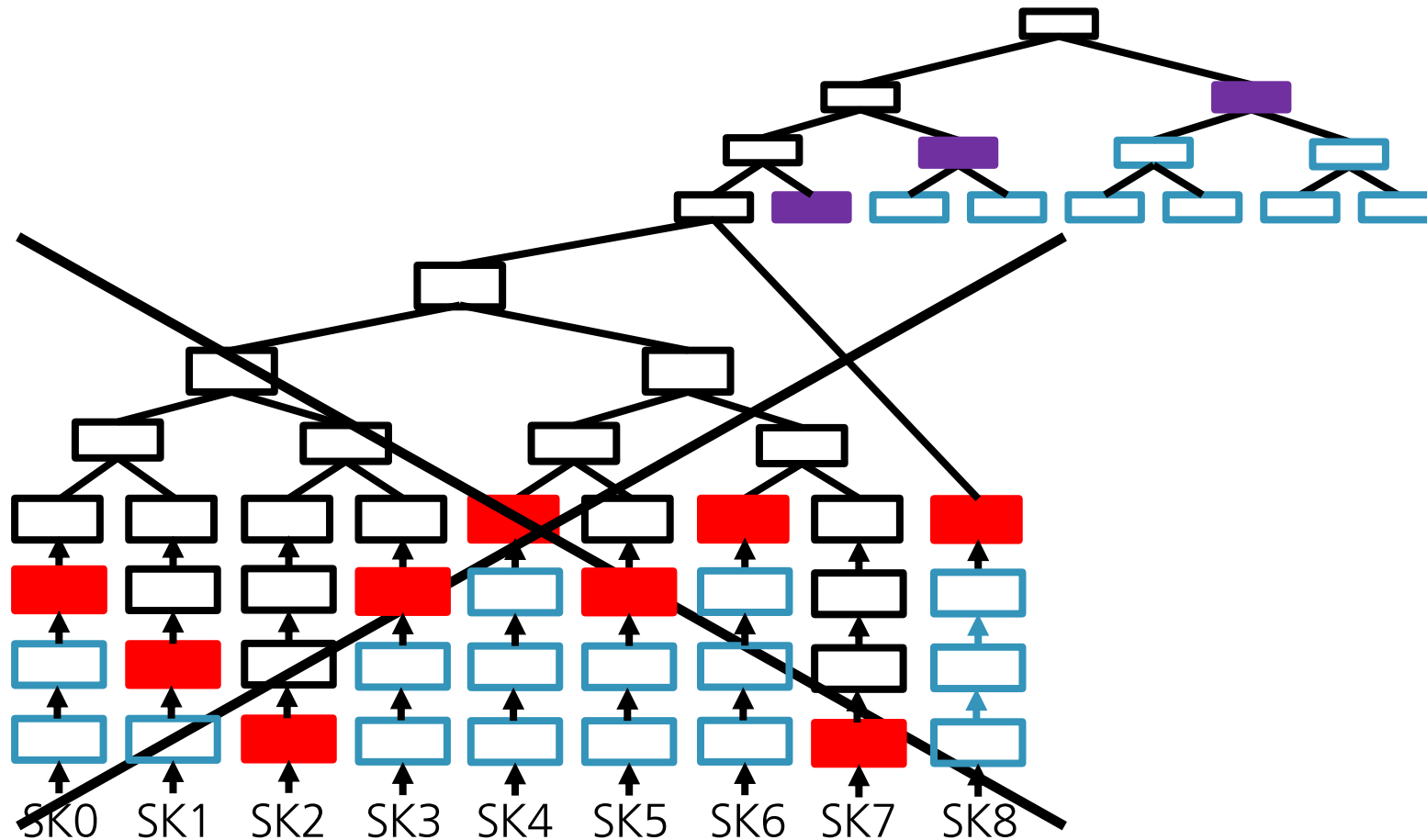
 Merkle Tree Authentication Path

 Node to be computed



# eXtended Merkle Signature Scheme

Signature Generation – Message 1




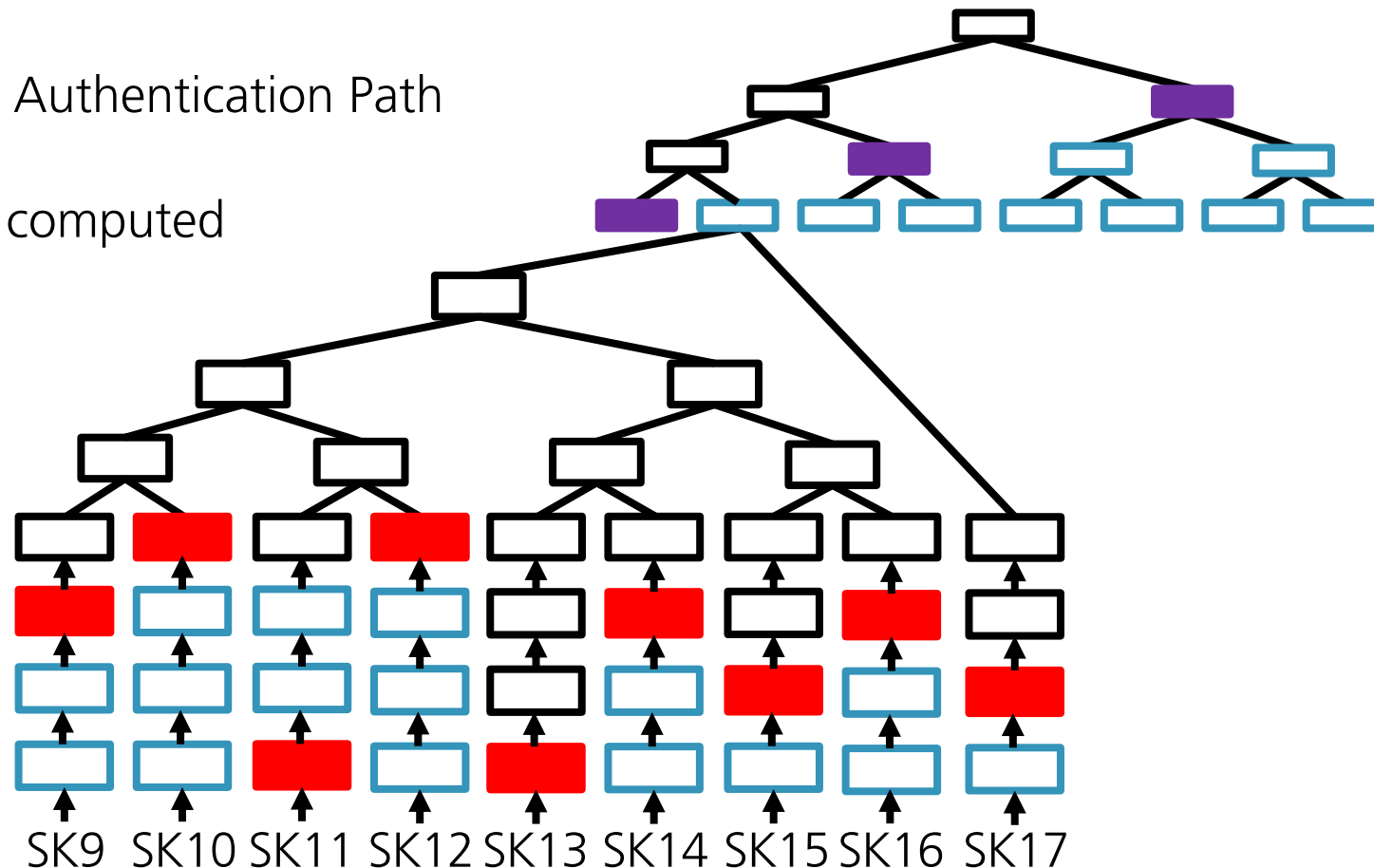
# eXtended Merkle Signature Scheme

Signature Generation – Message 2

 WOTS+ Signature




 Merkle Tree Authentication Path

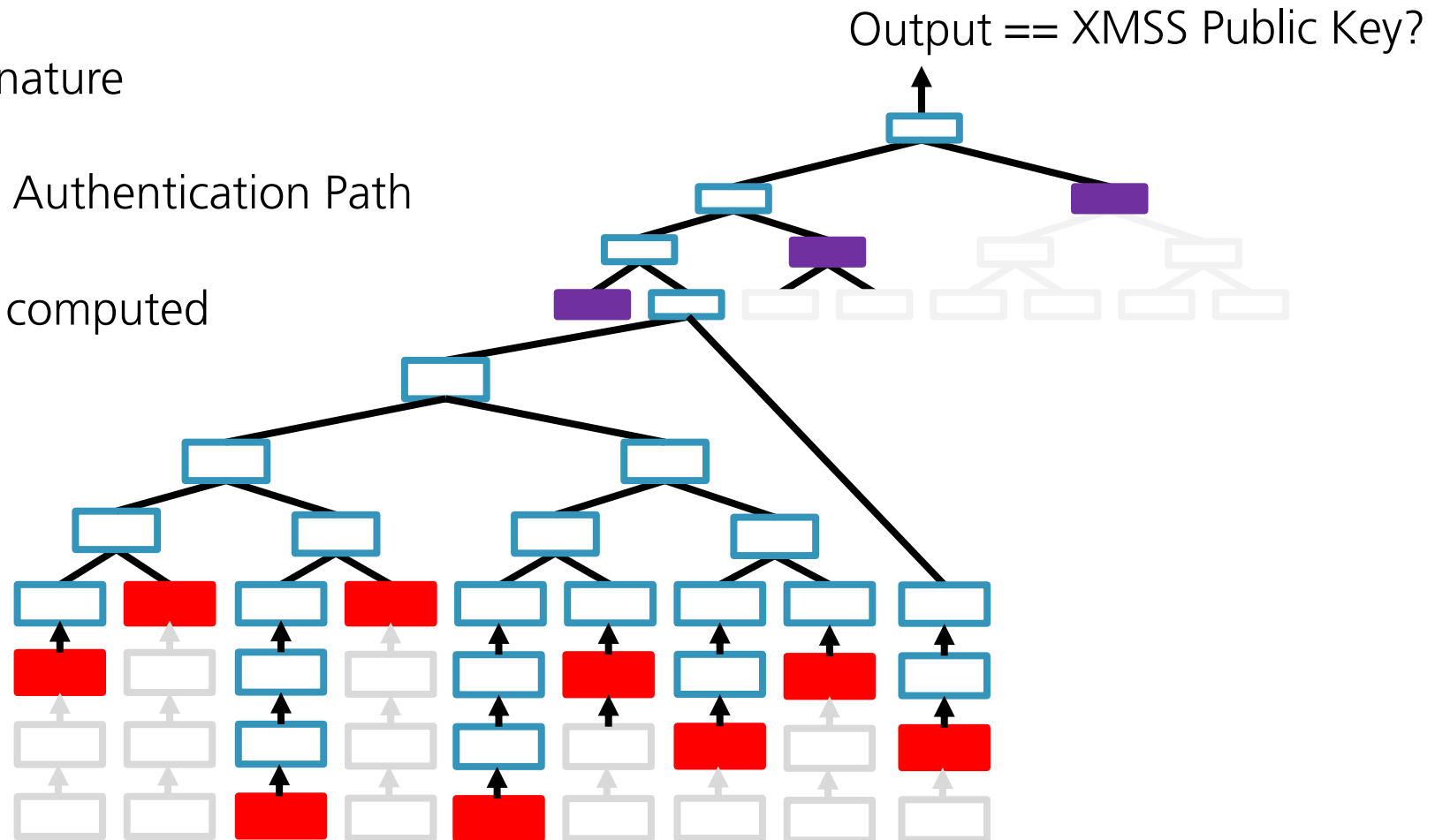
 Node to be computed



# eXtended Merkle Signature Scheme

Signature Verification – Message 2

-  WOTS+ Signature
-  Merkle Tree Authentication Path
-  Node to be computed



# Performance Estimates

# Performance Consideration

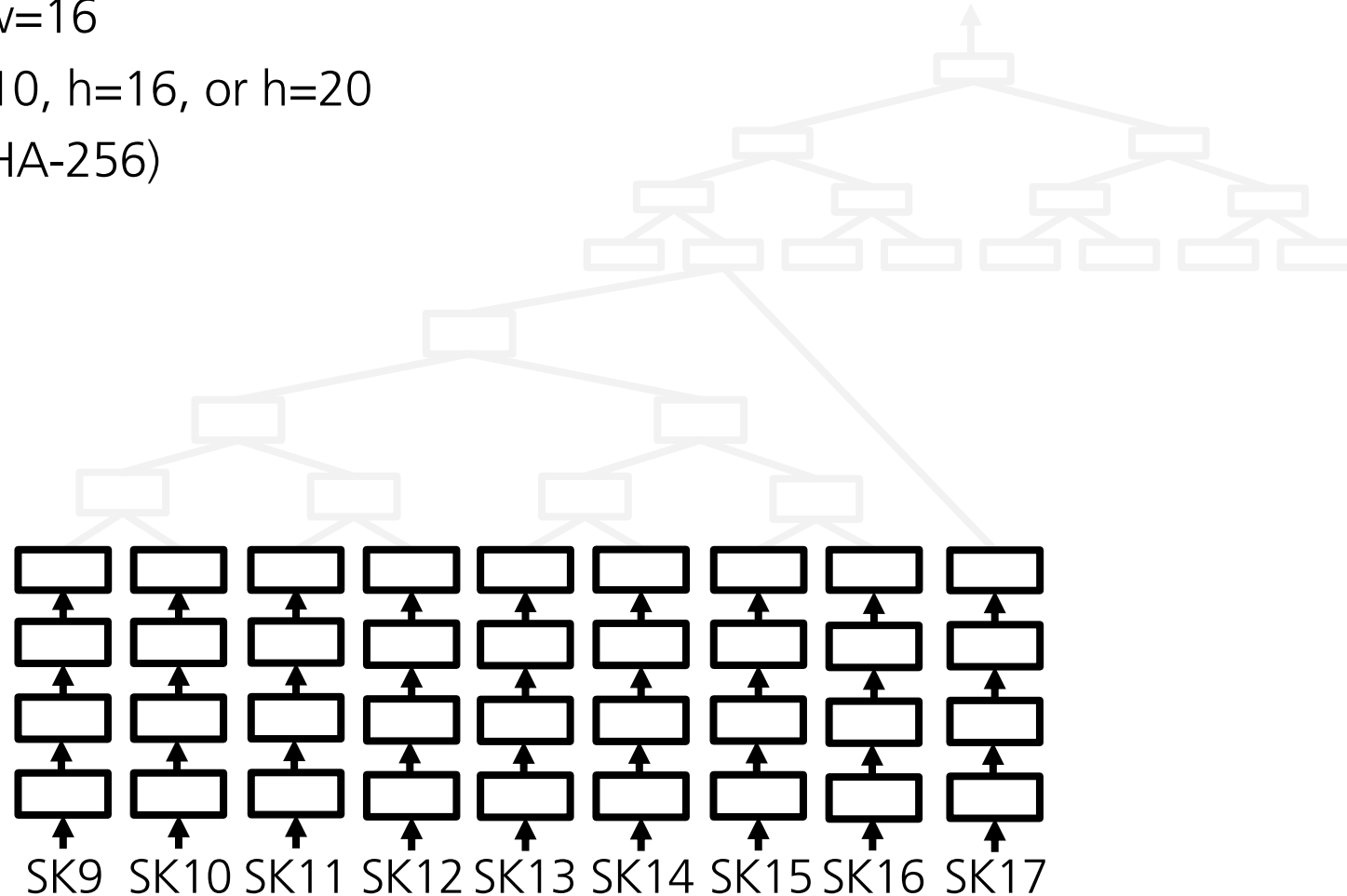
Public Key Generation – WOTS+

IRTF Parameters:

WOTS+ chain length  $w=16$

Merkle tree height  $h=10$ ,  $h=16$ , or  $h=20$

256 Bit Hashes (e.g. SHA-256)





# Performance Consideration

Public Key Generation – WOTS+

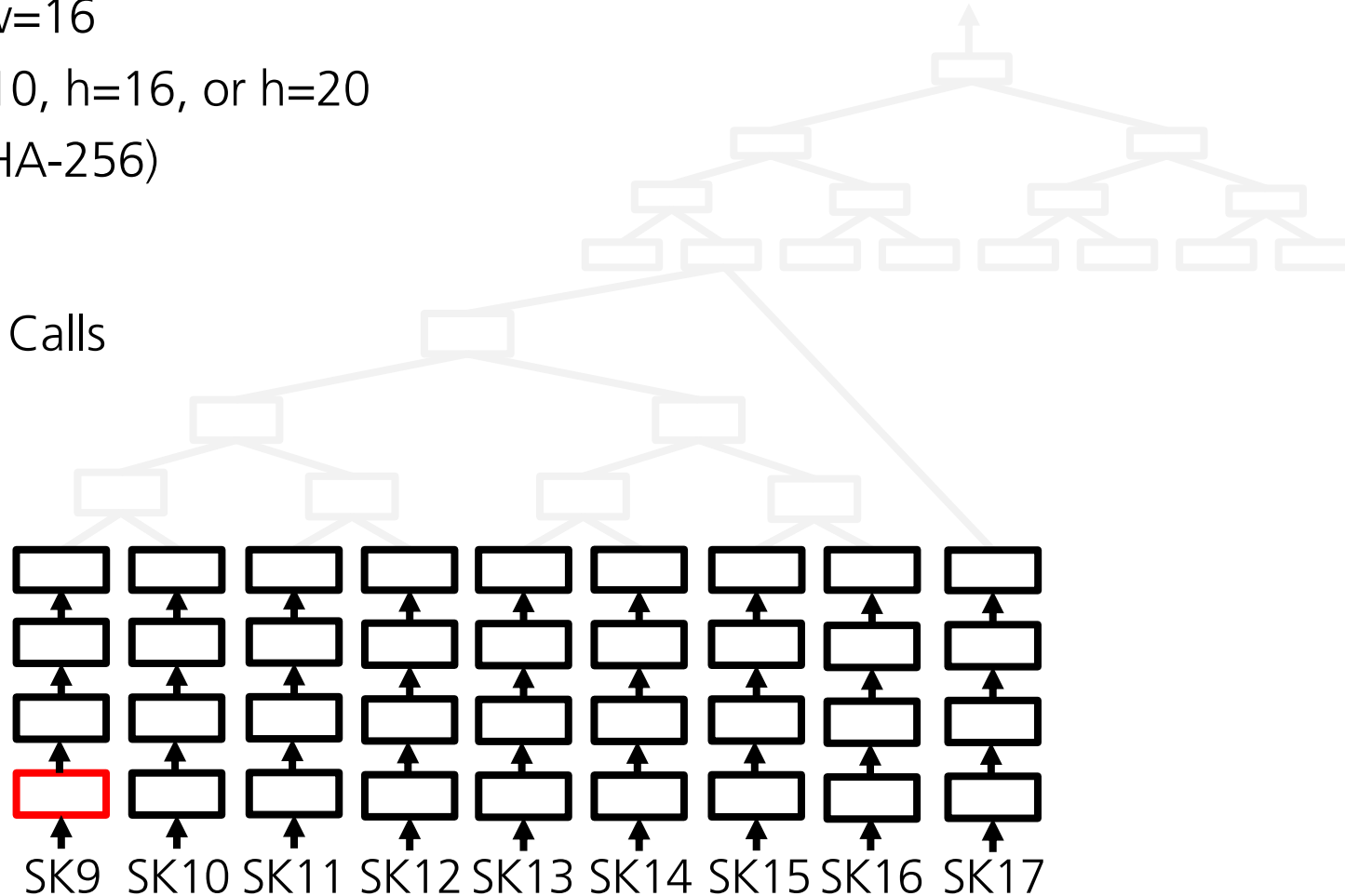
IRTF Parameters:

WOTS+ chain length  $w=16$

Merkle tree height  $h=10$ ,  $h=16$ , or  $h=20$

256 Bit Hashes (e.g. SHA-256)

 3 Hash Function Calls



# Performance Consideration


Public Key Generation – WOTS+

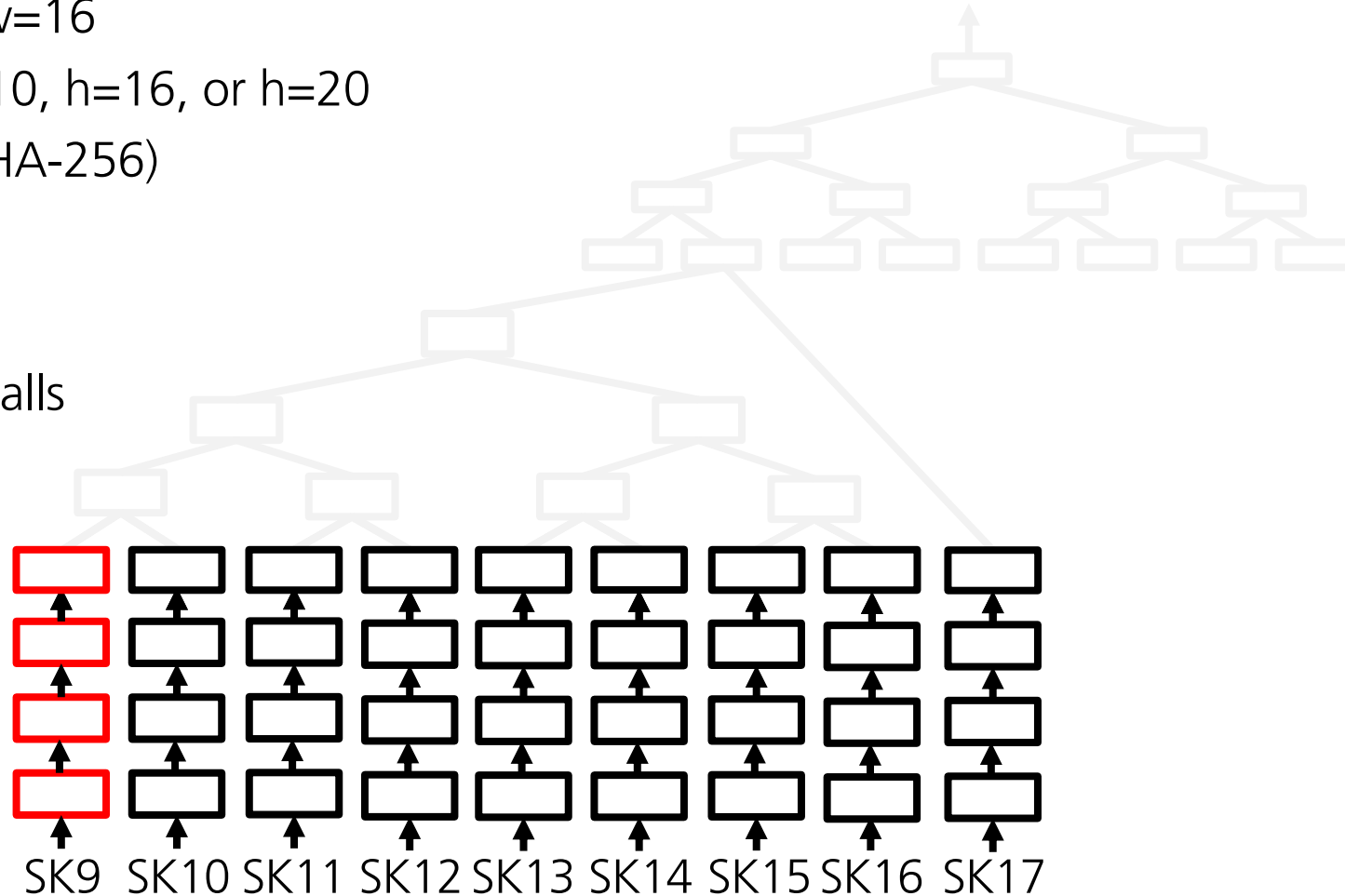
IRTF Parameters:

WOTS+ chain length  $w=16$

Merkle tree height  $h=10$ ,  $h=16$ , or  $h=20$

256 Bit Hashes (e.g. SHA-256)

  $3*w = 48$   
Hash Function Calls



# Performance Consideration


Public Key Generation – WOTS+

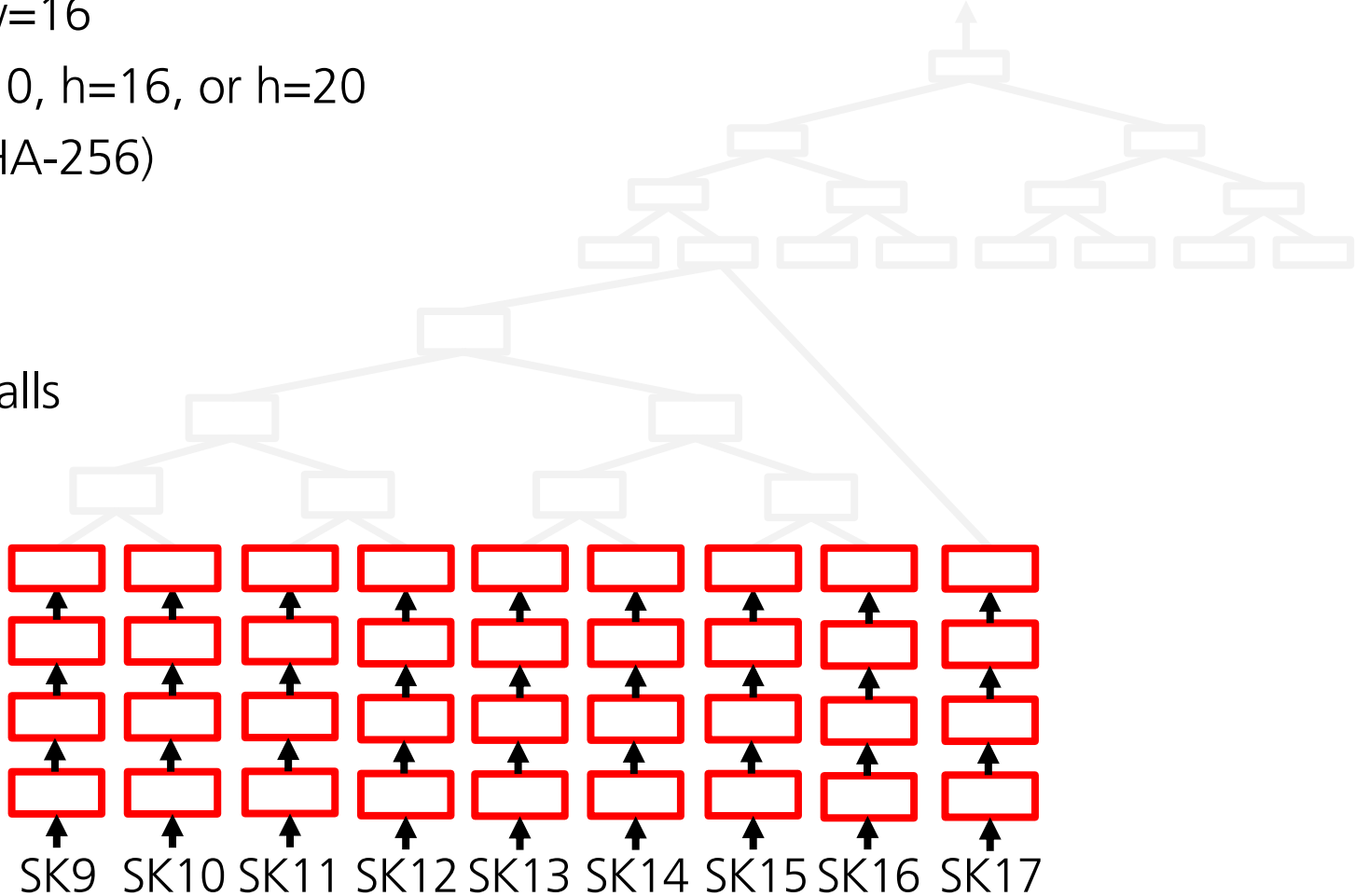
IRTF Parameters:

WOTS+ chain length  $w=16$

Merkle tree height  $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

  $48 \cdot 67 = 3216$   
Hash Function Calls



# Performance Consideration


Public Key Generation – WOTS+

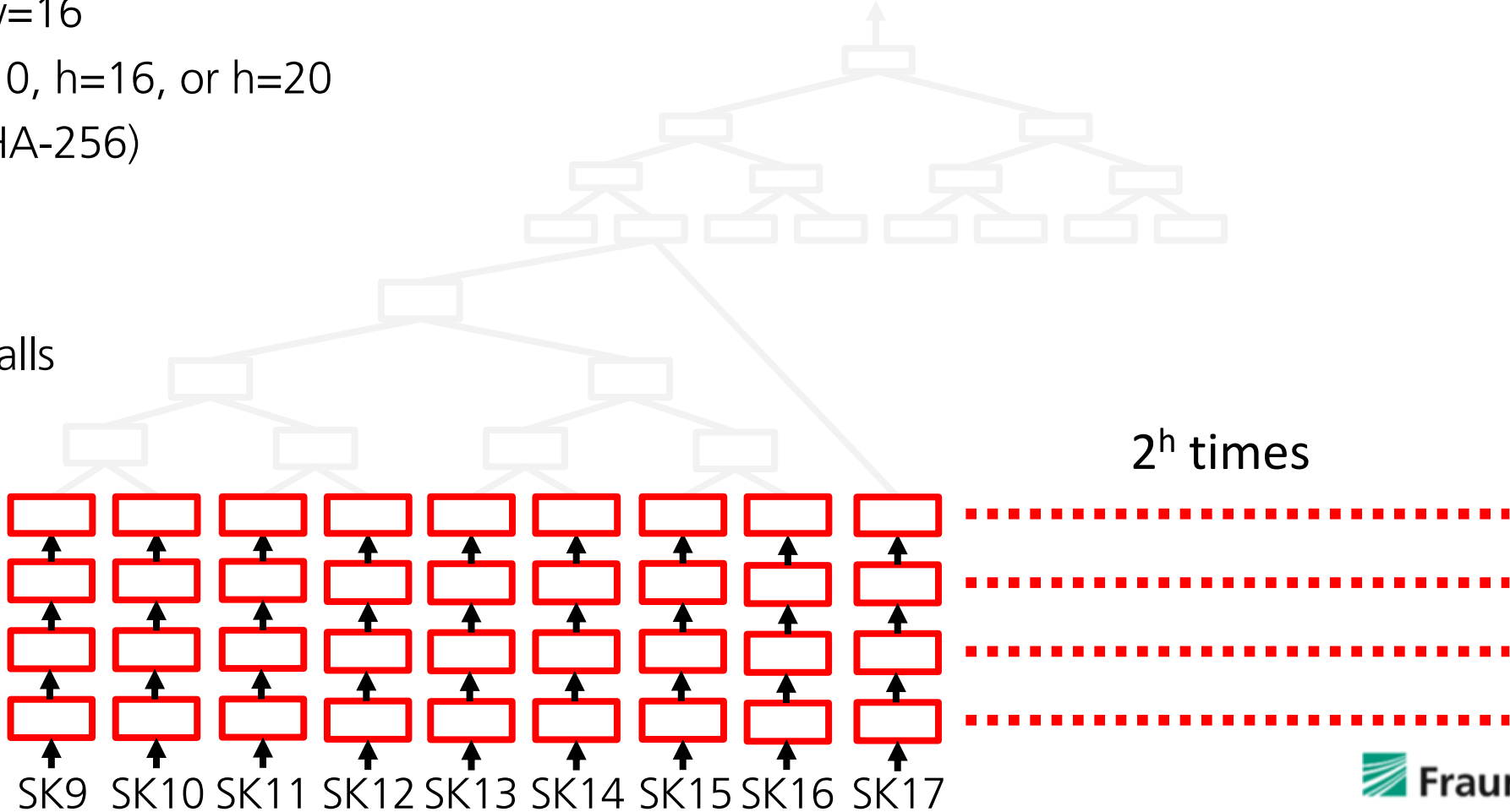
IRTF Parameters:

WOTS+ chain length  $w=16$

Merkle tree height  $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

  $3216 \cdot 2^h$   
Hash Function Calls



# Performance Consideration

Public Key Generation – L-Tree

IRTF Parameters:

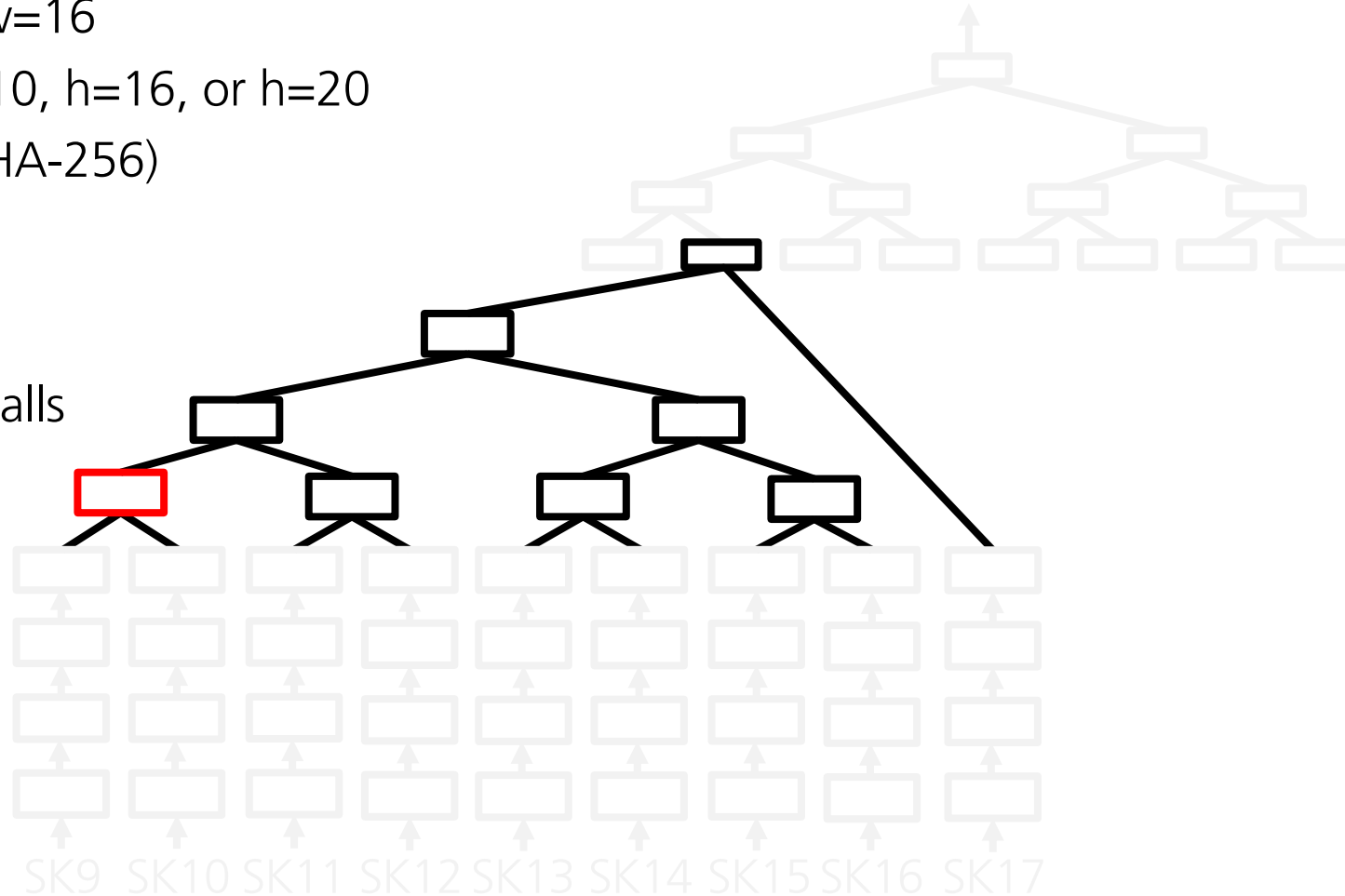
WOTS+ chain length  $w=16$

Merkle tree height  $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

 4

Hash Function Calls



# Performance Consideration


Public Key Generation – L-Tree

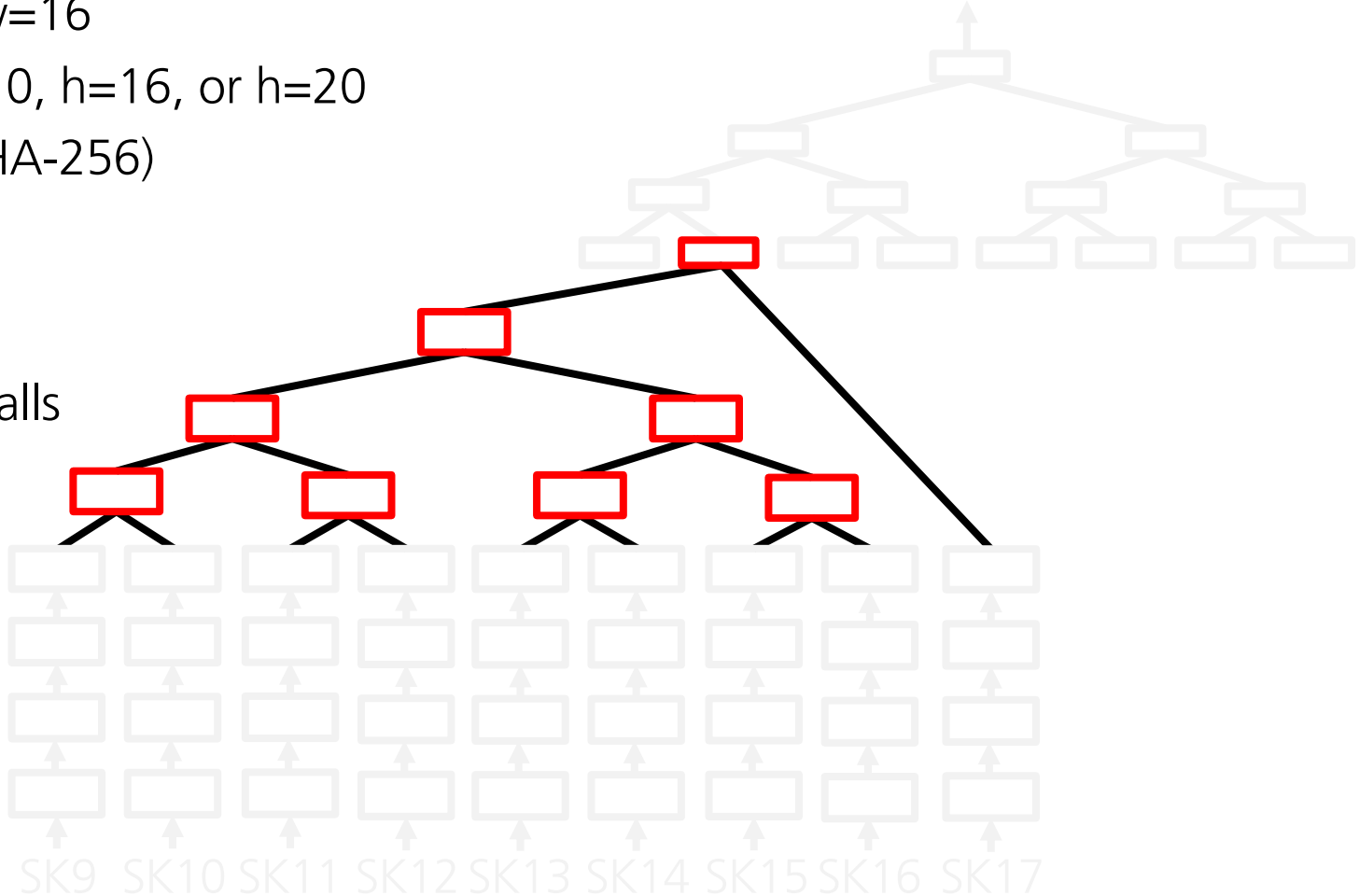
IRTF Parameters:

WOTS+ chain length  $w=16$

Merkle tree height  $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

  $4 \cdot 65 = 268$   
Hash Function Calls



# Performance Consideration


Public Key Generation – L-Tree

IRTF Parameters:

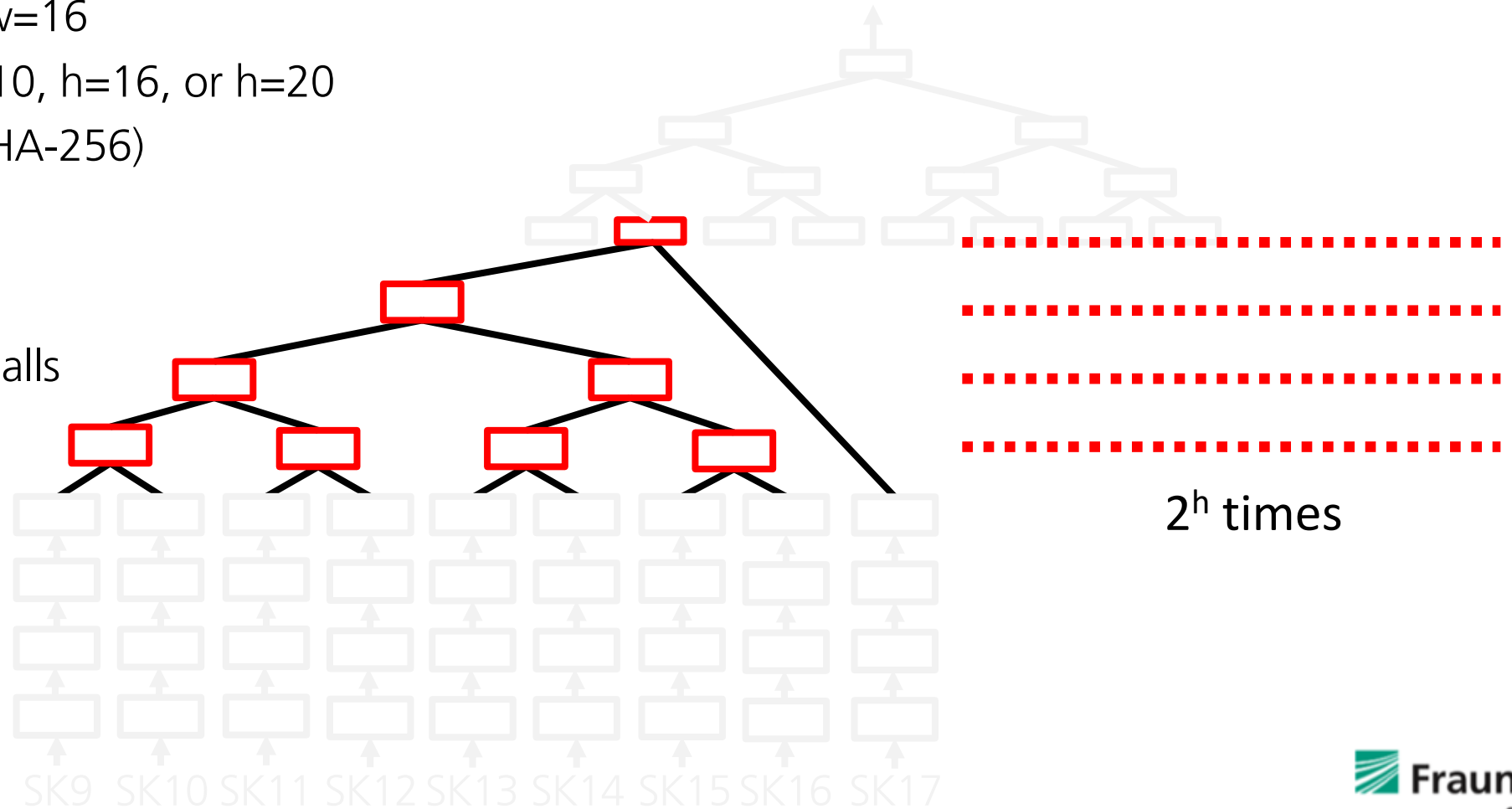
WOTS+ chain length  $w=16$

Merkle tree height  $h=10$ ,  $h=16$ , or  $h=20$

256 Bit Hashes (e.g. SHA-256)

  $260 \cdot 2^h$

Hash Function Calls



# Performance Consideration


Public Key Generation – XMSS

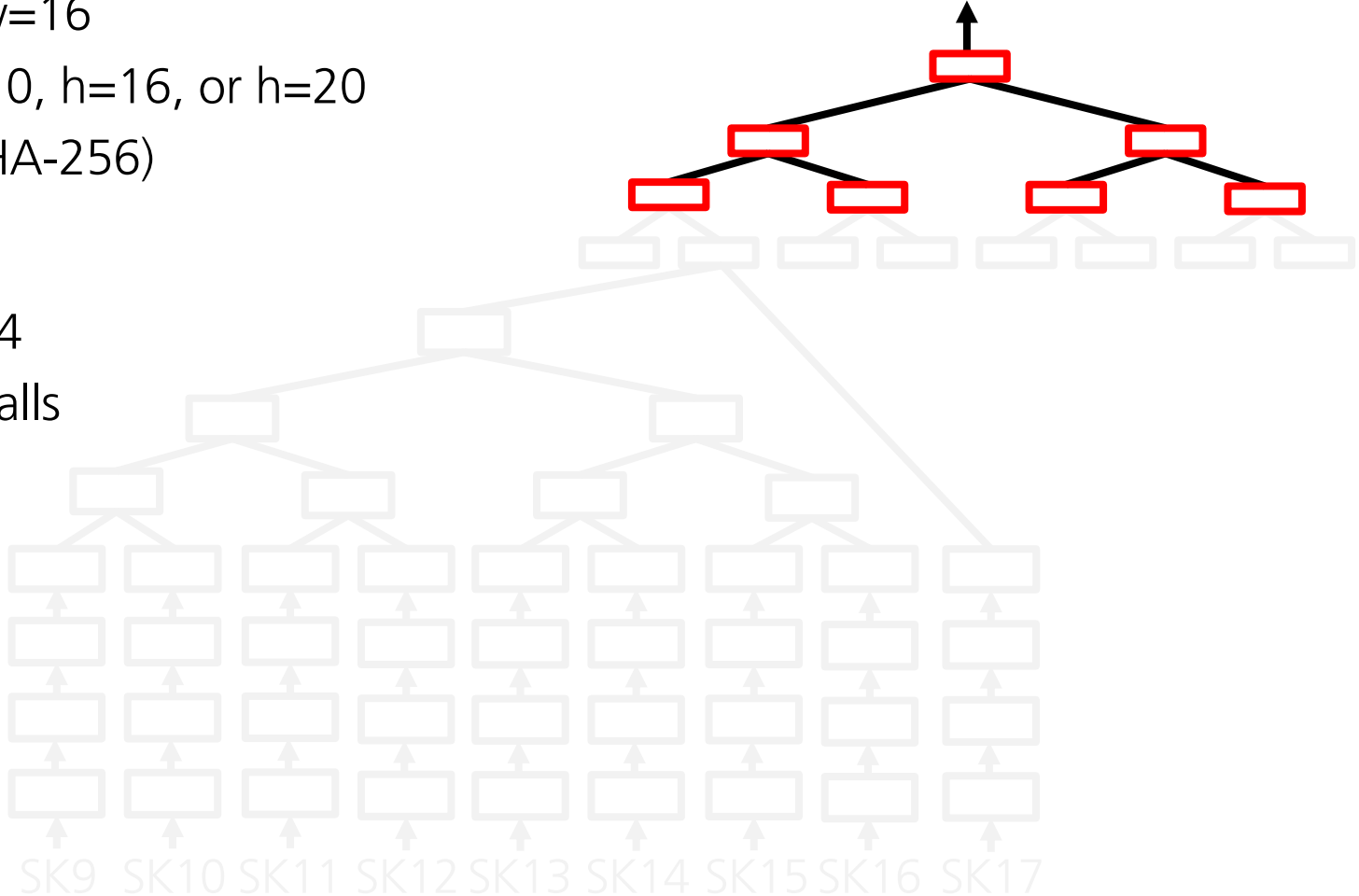
IRTF Parameters:

WOTS+ chain length  $w=16$

Merkle tree height  $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

  $4 \cdot (2^h - 1) = 4 \cdot 2^h - 4$   
Hash Function Calls





# Performance Consideration


Public Key Generation – XMSS

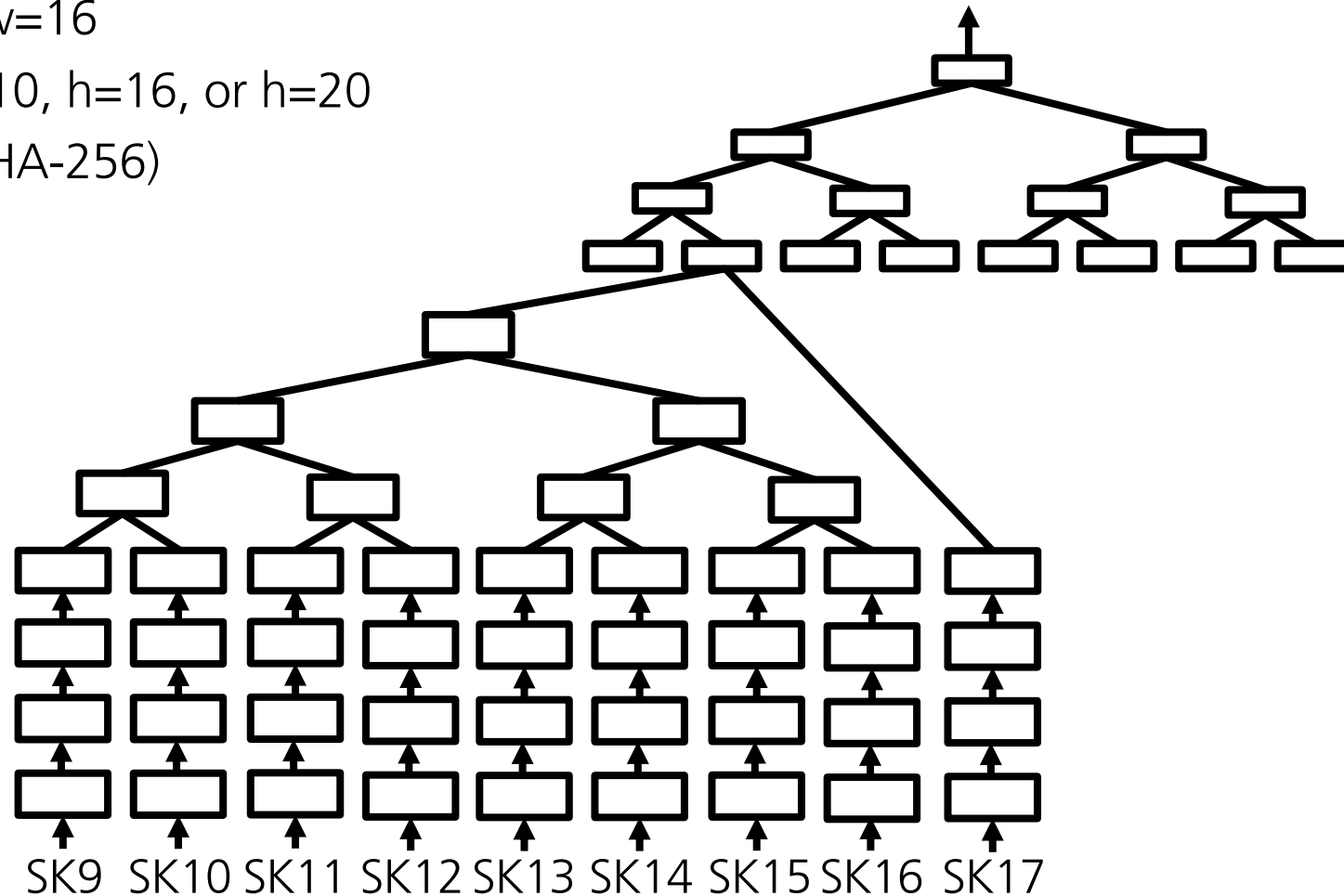
IRTF Parameters:

WOTS+ chain length  $w=16$

Merkle tree height  $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

  $3480 * 2^h - 4$   
Total Hash  
Function Calls



# Performance Consideration

## Hash Function Calls

|                        | h=10      | h=16        | h=20          |
|------------------------|-----------|-------------|---------------|
| Signatures             | 1024      | 65,536      | 1,048,576     |
| Public Key Generation  | 3,563,520 | 228,065,280 | 3,649,044,480 |
| Signature Generation   | ~5,560    | ~263,684    | ~4,195,828    |
| Signature Verification | ~1,908    | ~1,932      | ~1,948        |

# Performance with SHA-256

|                       | h=10                     | h=16                         | h=20                          |
|-----------------------|--------------------------|------------------------------|-------------------------------|
| Signatures            | 1024                     | 65,536                       | 1,048,576                     |
| Public Key Generation | 423,099,648 clock cycles | $27 \cdot 10^9$ clock cycles | $434 \cdot 10^9$ clock cycles |
| With 400 MHz          | <1.1 s                   | <70 s                        | <1085 s                       |
| Sign                  | < 2 ms                   | < 70 ms                      | < 1 s                         |
| Verify                | < 1 ms                   | < 1 ms                       | < 1 ms                        |

# Performance with SHA-3

|                       | h=10                    | h=16                    | h=20                     |
|-----------------------|-------------------------|-------------------------|--------------------------|
| Signatures            | 1024                    | 65,536                  | 1,048,576                |
| Public Key Generation | 79,159,200 clock cycles | $5 * 10^9$ clock cycles | $81 * 10^9$ clock cycles |
| With 400 MHz          | < 200 ms                | < 12.5 s                | < 203 s                  |
| Sign                  | < 1 ms                  | < 12.5 ms               | < 200 ms                 |
| Verify                | < 1 ms                  | < 1 ms                  | < 1 ms                   |

# Comparison with ECC

FPGA Implementation Estimates (Virtex-5)

|                       | Ed25519 | XMSS-SHA3 h=10 |
|-----------------------|---------|----------------|
| Public Key Generation | < 1 ms  | < 200 ms       |
| Sign                  | < 1 ms  | < 1 ms         |
| Verify                | < 2 ms  | < 1 ms         |

# Optimisations and Trade-Offs

## Parallelization and Caching

- Parallelization
  - WOTS+ trivial to compute in parallel
  - L-Tree and XMSS more difficult to parallelize
- More/Less Caching
  - More caching of XMSS for authentication path (costs more memory)
    - ➔ Improves the signing performance
  - Less caching to save memory
    - ➔ In the worst case, signing almost as slow as public key generation
    - ➔ Useful for lightweight applications with low memory



Thank you for your attention!