

Hash-based Signatures

IETF/IRTF CFRG Draft on XMSS



Fraunhofer Workshop Series 01 – Post-Quantum Cryptography in Practice
Speaker: Dr. Bernhard Jungk

eXtended Merkle Signature Scheme

eXtended Merkle Signature Scheme

Why should we look into XMSS?

Hash-based signatures have many advantages:

- Based on well understood security notions
 - » Cryptographic hash functions are hard to invert, **also for quantum computers**
 - » Merkle trees well studied since the 1980ies
- Hash functions are well understood (especially after SHA-3 competition)
- Fast signing and verification operations possible
- Relatively easy to understand and implement

eXtended Merkle Signature Scheme

Why should we look into XMSS?

XMSS is a promising candidate for

- Applications with relatively low amount of signatures
- One- or many-times firmware updates
- Digital signatures for documents (e.g. contracts, email)
- Long-term archival of important digital assets
- PKI Certificates (e.g. Root CA)

eXtended Merkle Signature Scheme

Why should we look into XMSS?

IRTF is part of IETF

- Oriented towards research and long-term trends

Important trend – PQC

- Quantum computer attacks are likely
- Design of replacements for traditional public key crypto

Standardization needed

- Interoperability
- Implementation Guidelines

eXtended Merkle Signature Scheme

Our Contribution

Implementation experience

- Benchmarking against other schemes
- Learn good trade-offs for different application scenarios, cost reductions, side-channels, etc.

Target Platform: Hardware, i.e. FPGAs and ASICs

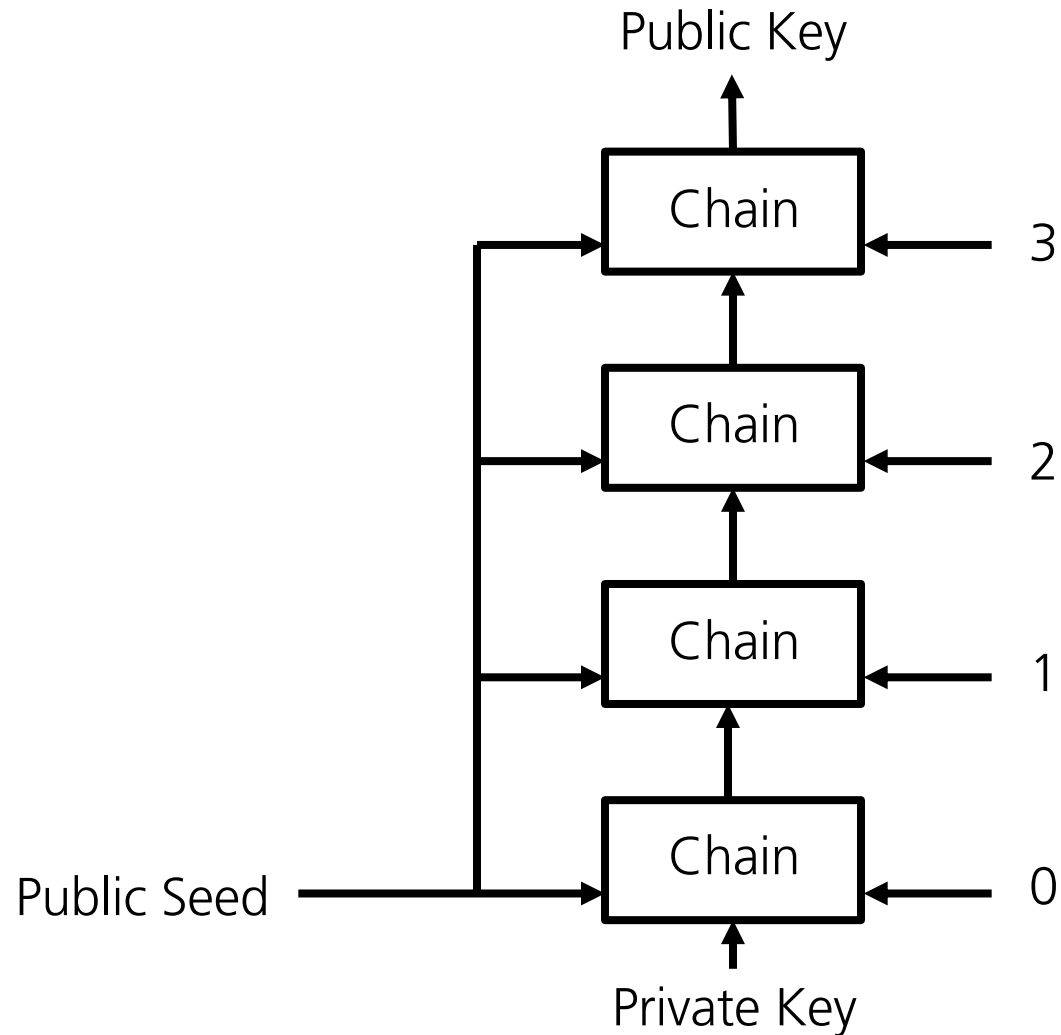
Cooperation:

- Yale University in New Haven, US
- Fraunhofer SIT in Darmstadt, Germany
- Fraunhofer Singapore

Recap Winternitz One-Time Signatures

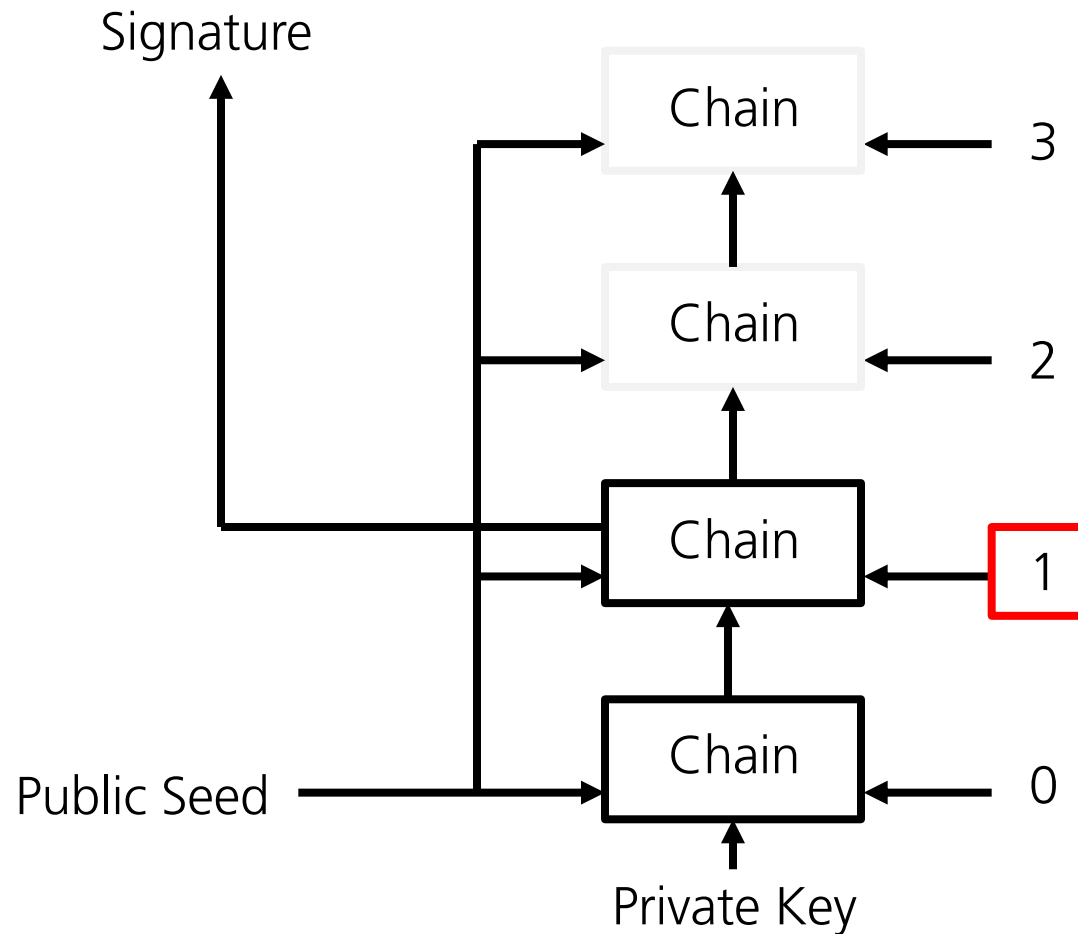
Winternitz One-Time Scheme+

Basic Principle – Public Key Generation



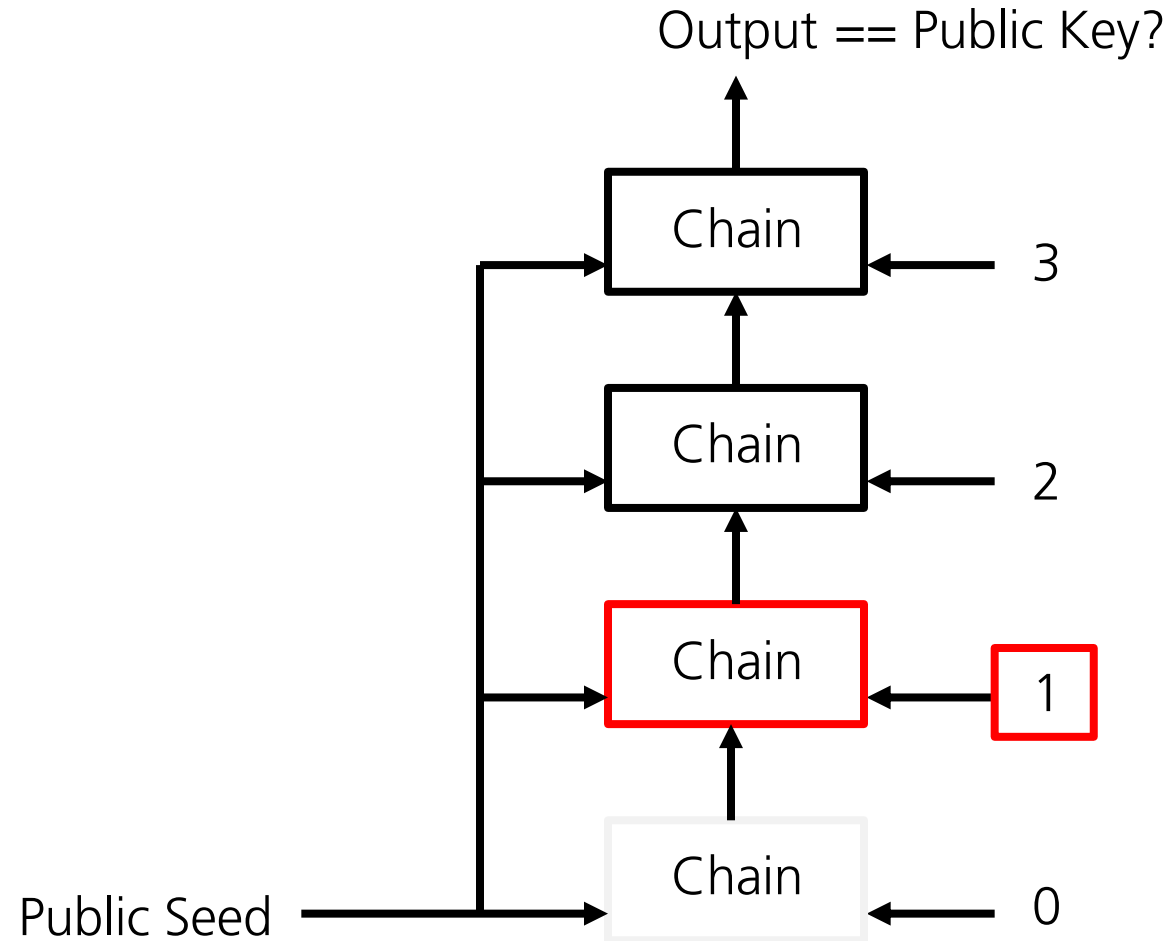
Winternitz One-Time Scheme+

Basic Principle – Signature Generation



Winternitz One-Time Scheme+

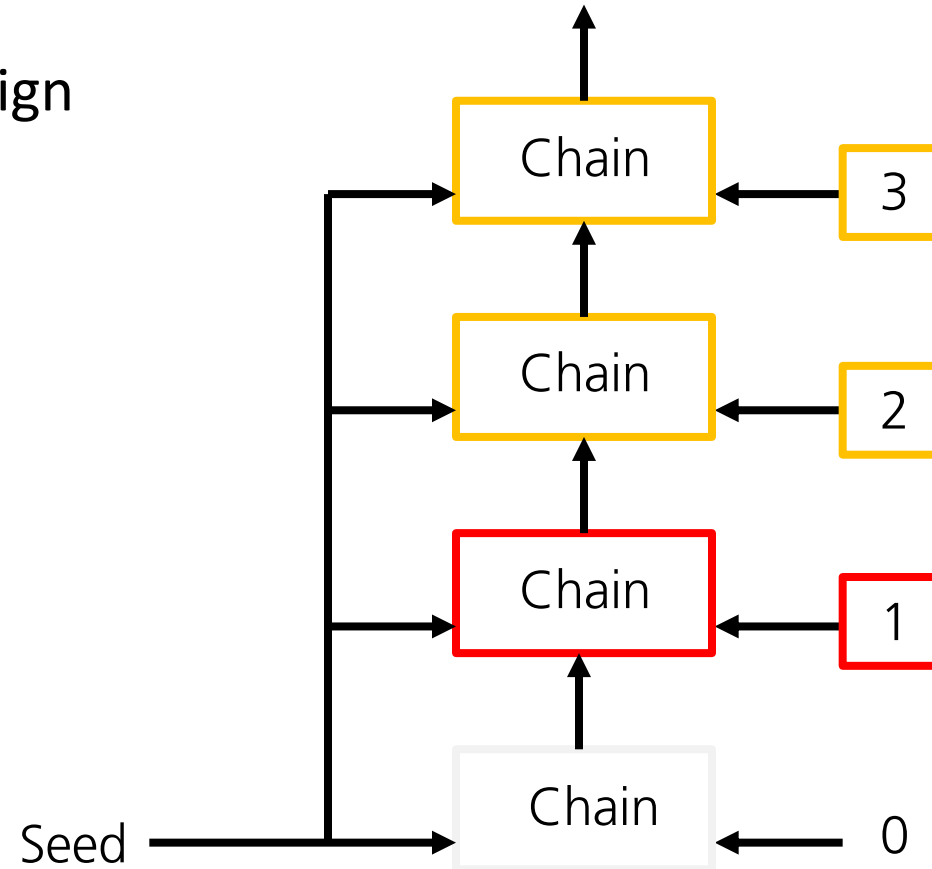
Basic Principle – Signature Verification



Winternitz One-Time Scheme+

Basic Principle

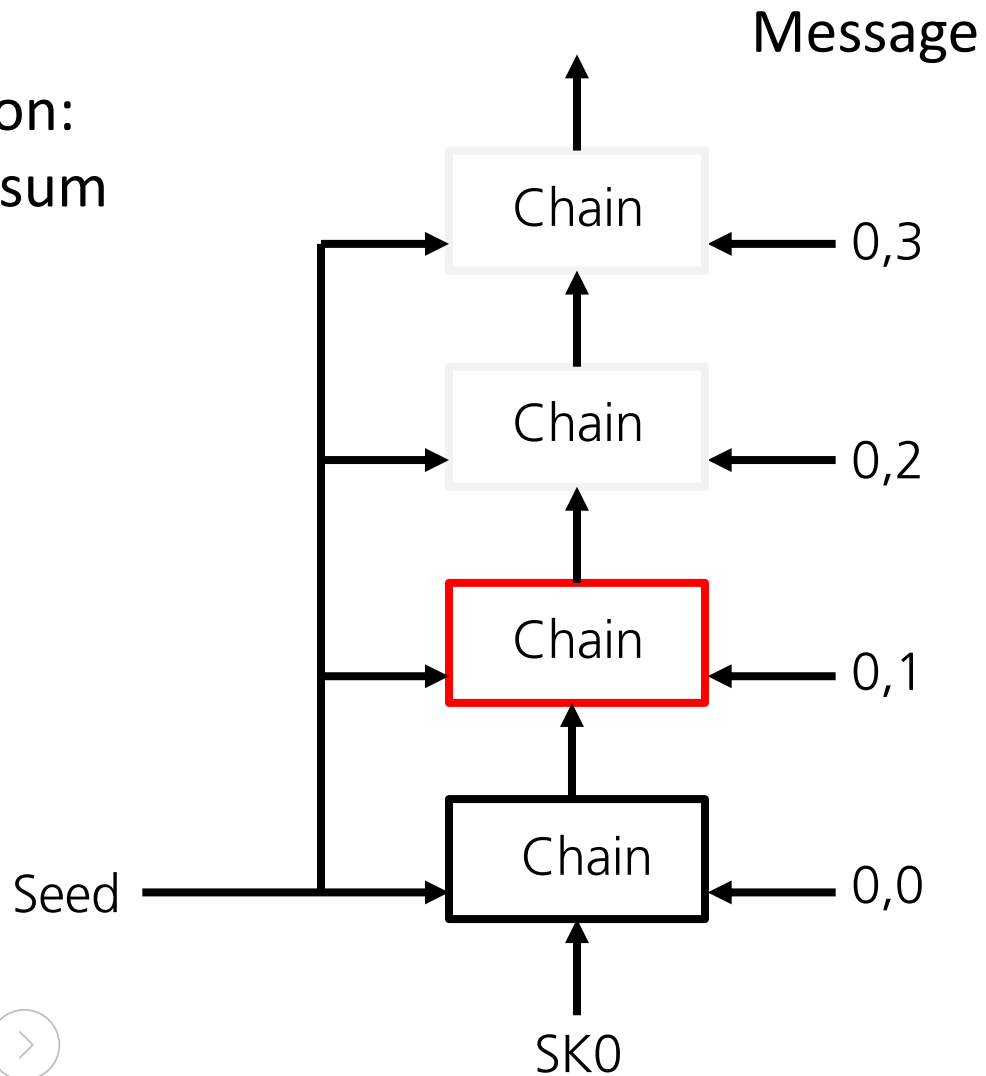
Problem:
Signer reveals how to sign
other messages with
the same key



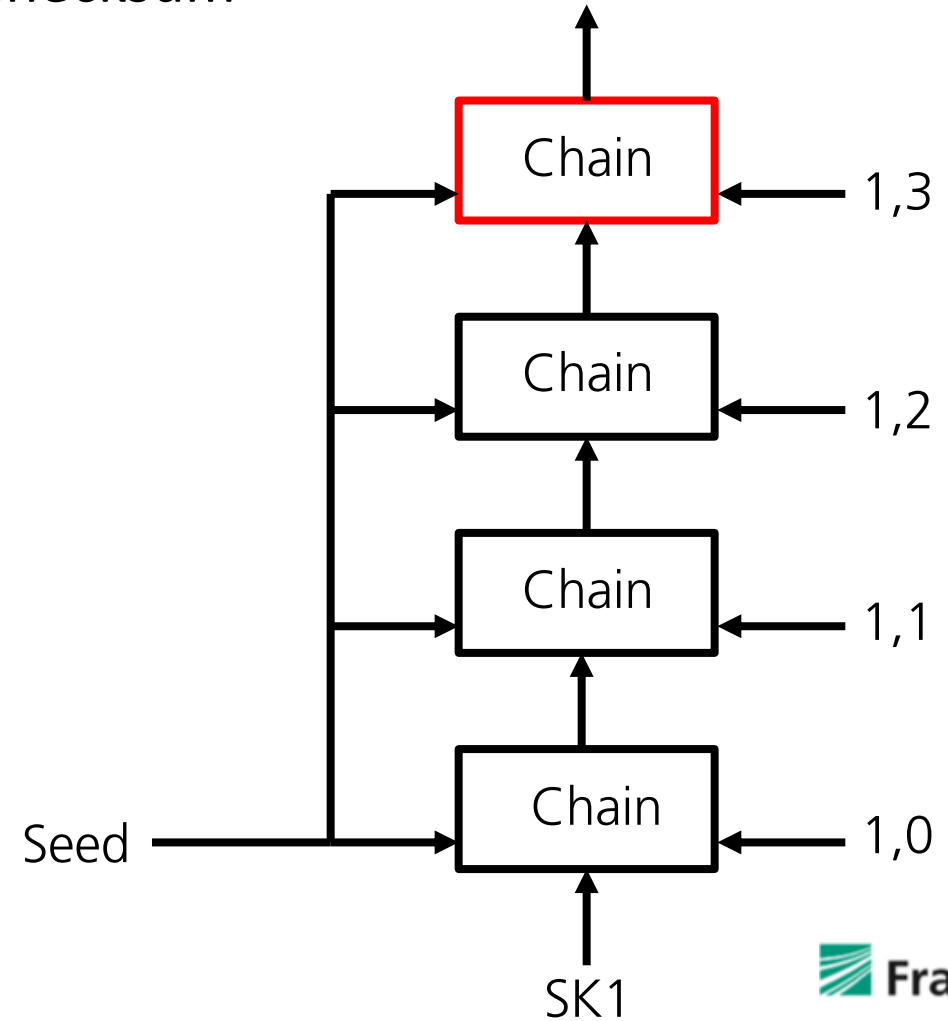
Winternitz One-Time Scheme+

Basic Principle

Solution:
Checksum

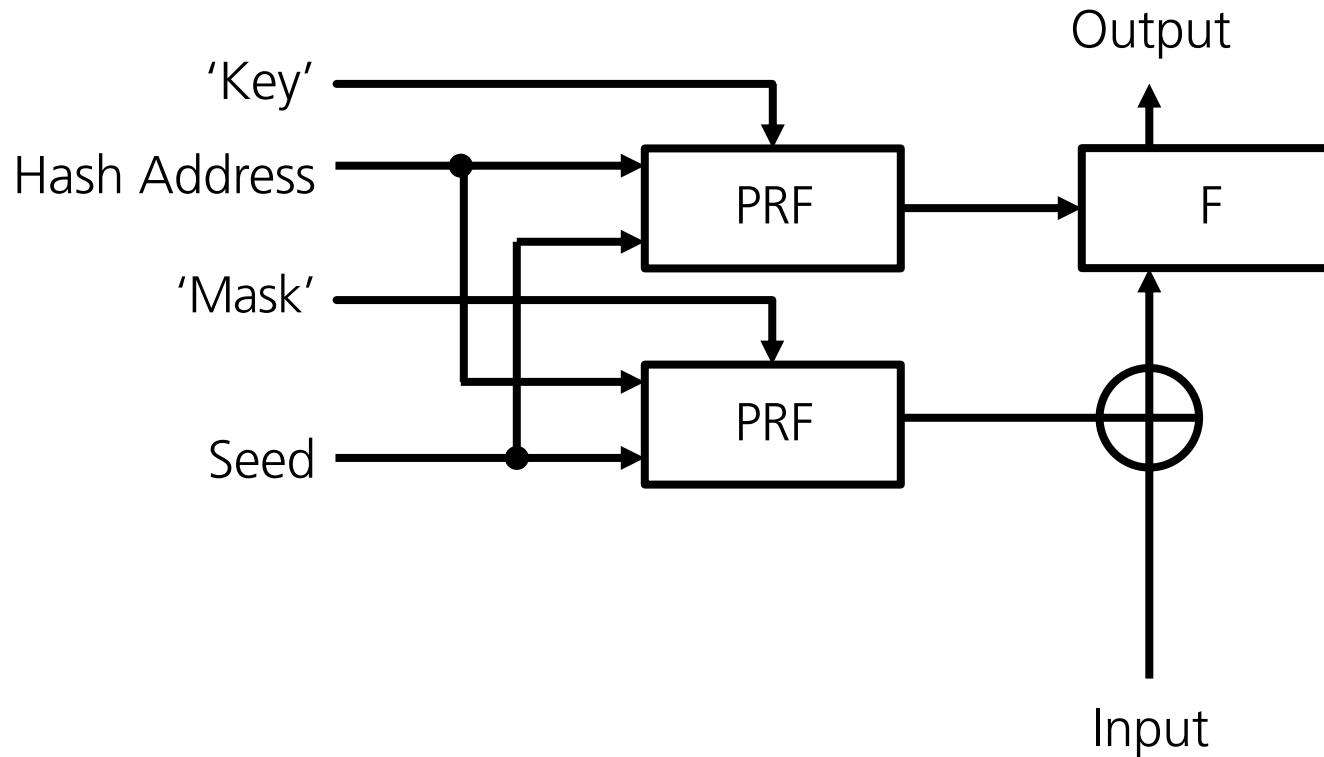


Checksum



Winternitz One-Time Scheme+

Chaining Function for XMSS

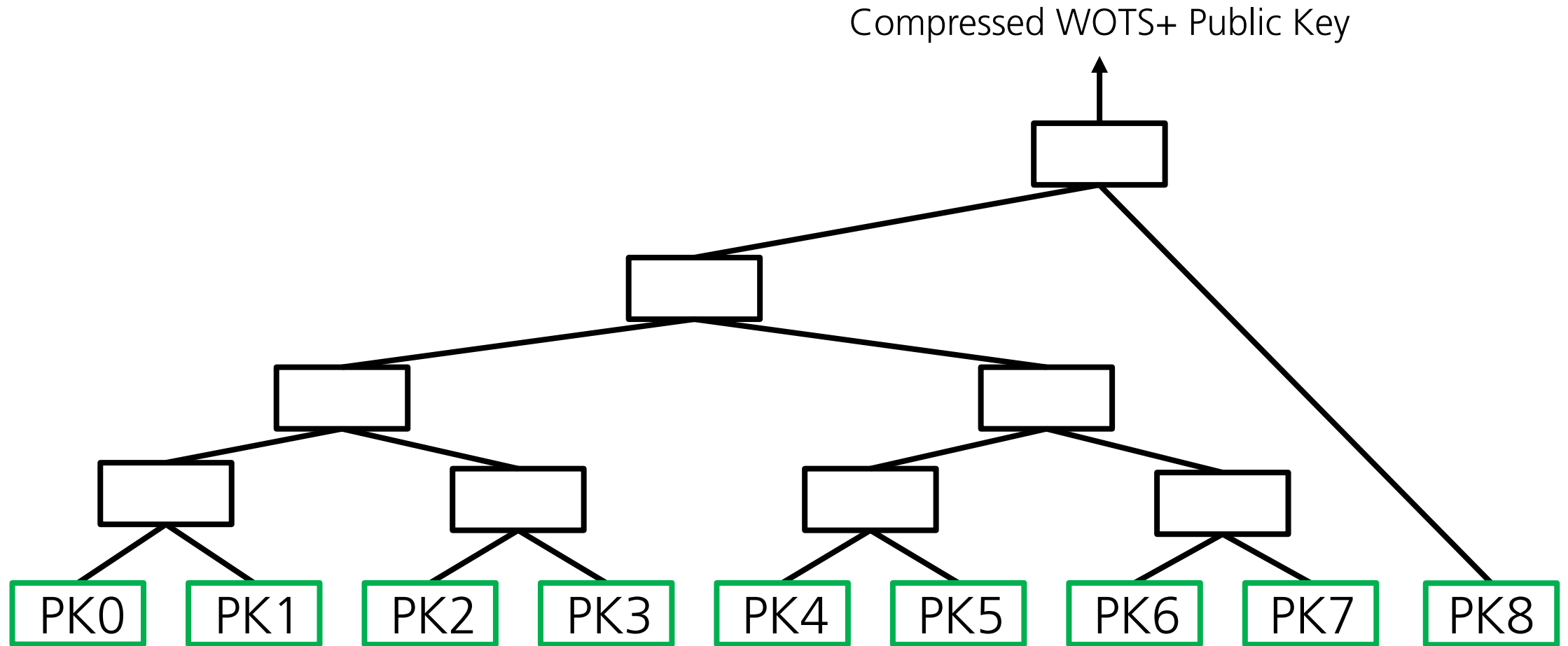


PRF – Pseudorandom function
F – Keyed hash function

eXtended Merkle Signature Scheme

eXtended Merkle Signature Scheme

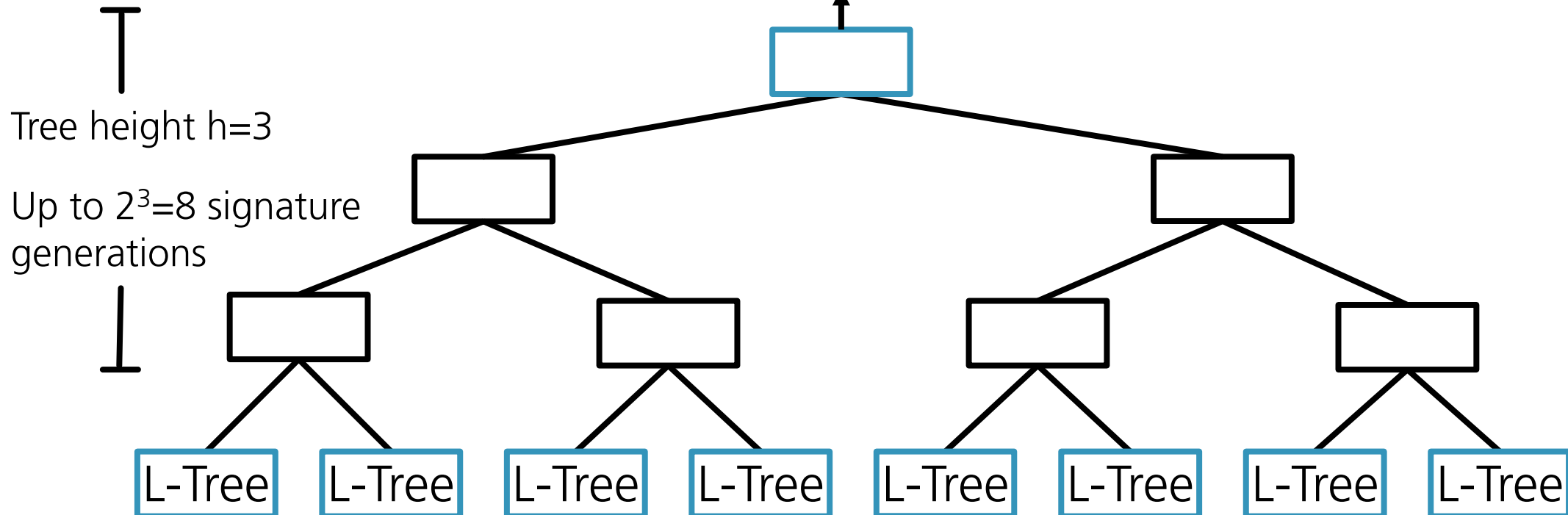
L-Tree – Public Key Generation



eXtended Merkle Signature Scheme

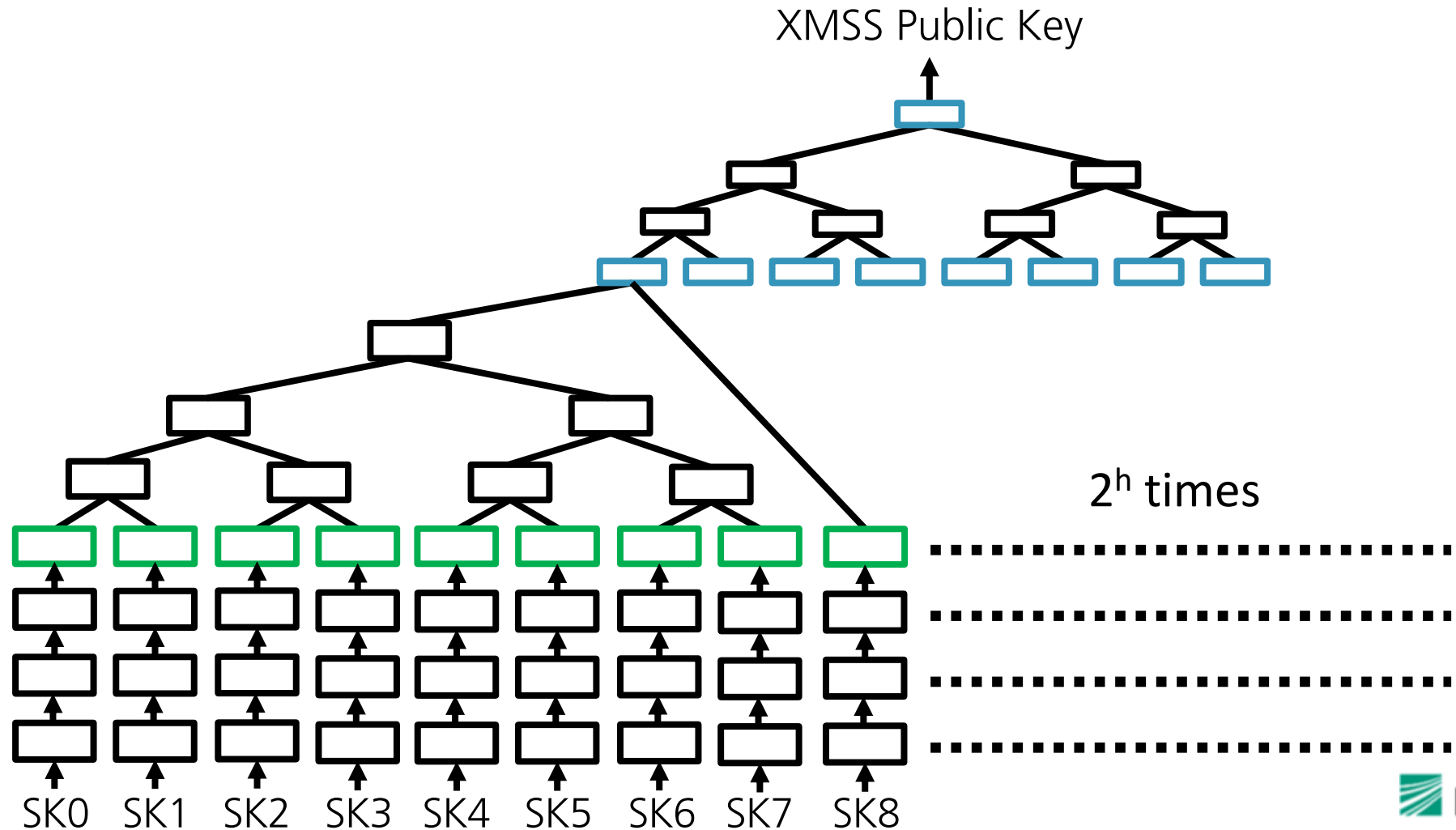
XMSS Tree – Public Key Generation

XMSS Public Key



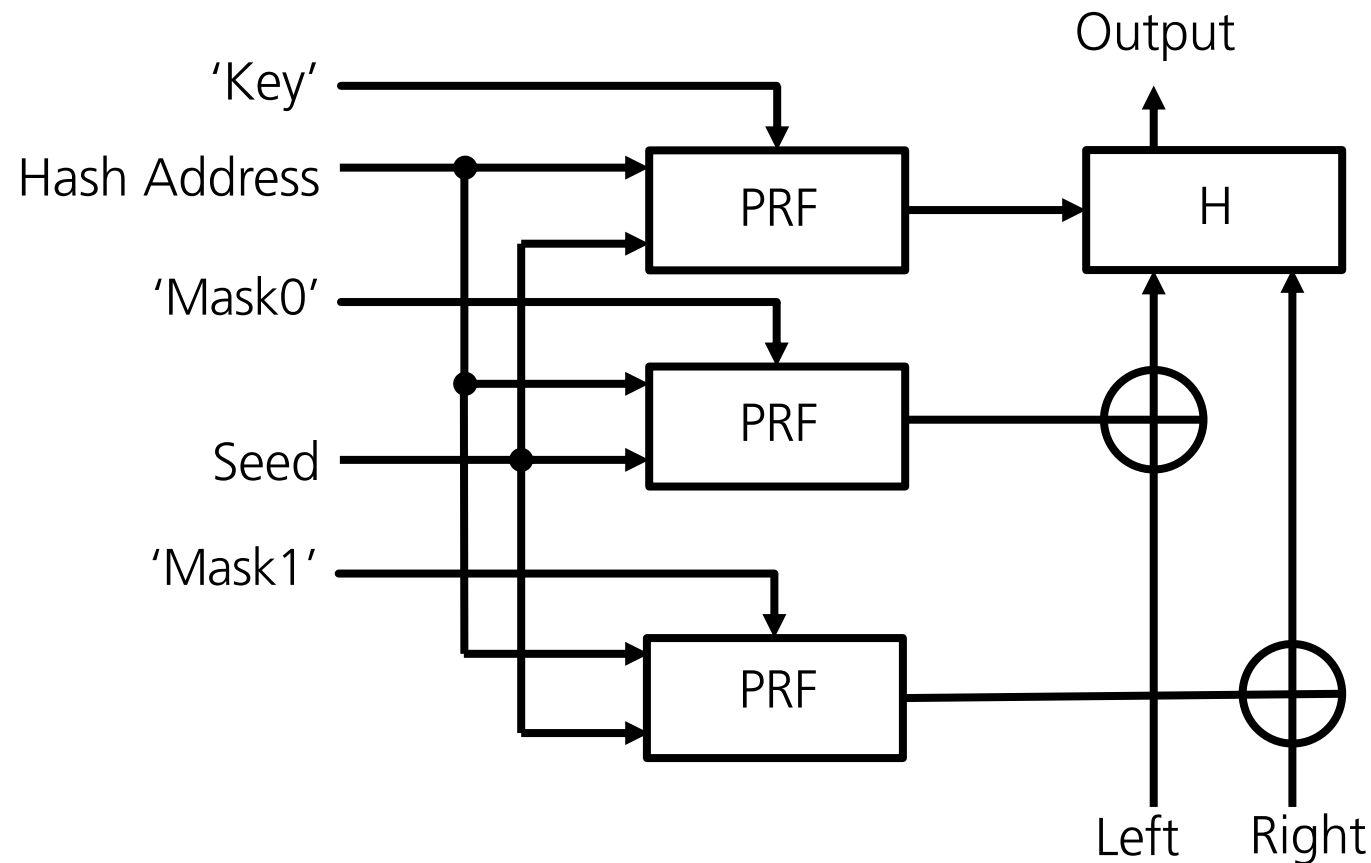
eXtended Merkle Signature Scheme

The Complete Picture – Public Key Generation



eXtended Merkle Signature Scheme

rand_hash



PRF – Pseudorandom function
H – Keyed hash function

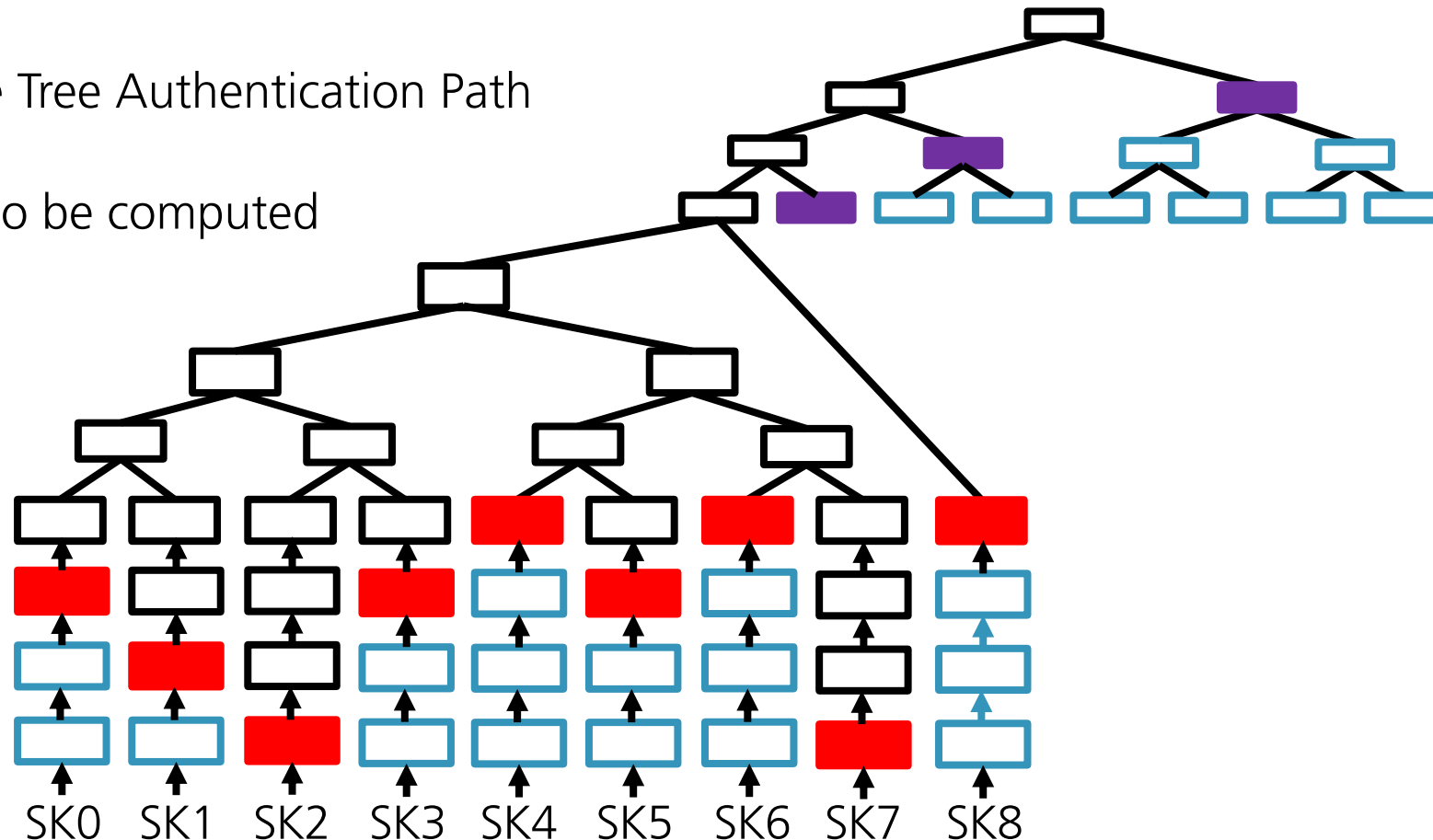
eXtended Merkle Signature Scheme

Signature Generation – Message 1

 WOTS+ Signature

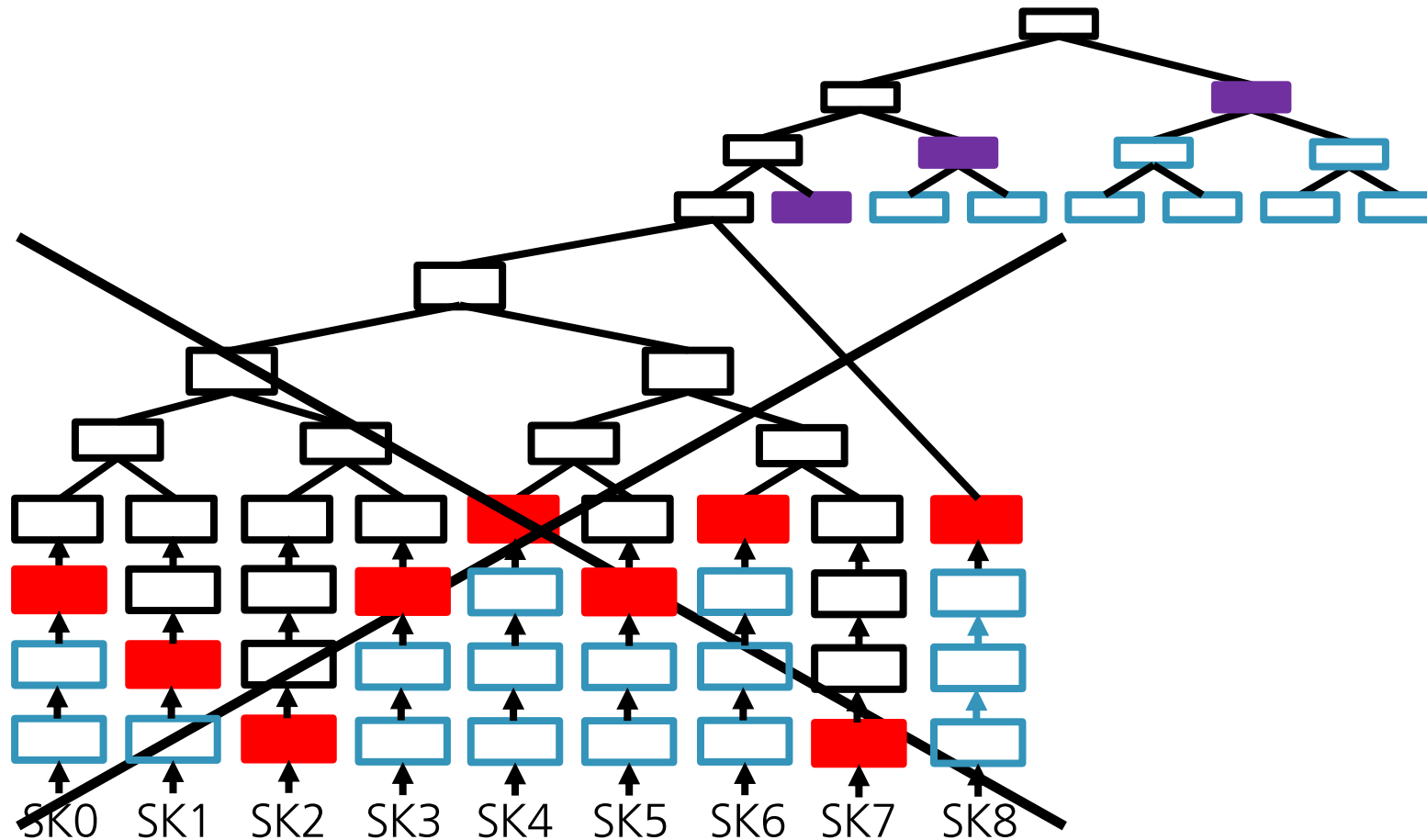
 Merkle Tree Authentication Path

 Node to be computed



eXtended Merkle Signature Scheme

Signature Generation – Message 1




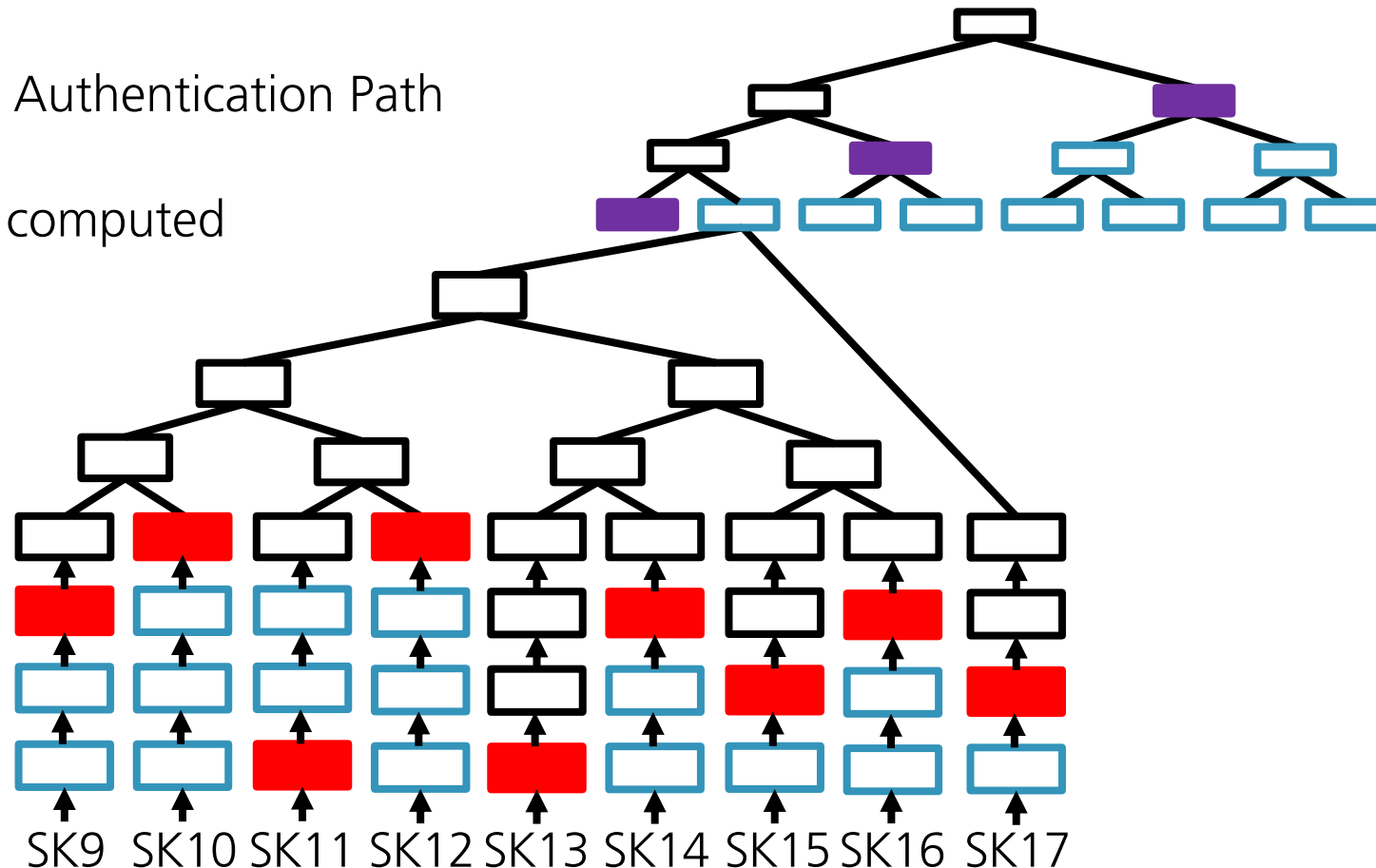
eXtended Merkle Signature Scheme

Signature Generation – Message 2

 WOTS+ Signature

 Merkle Tree Authentication Path

 Node to be computed



Performance Estimates

Performance Consideration

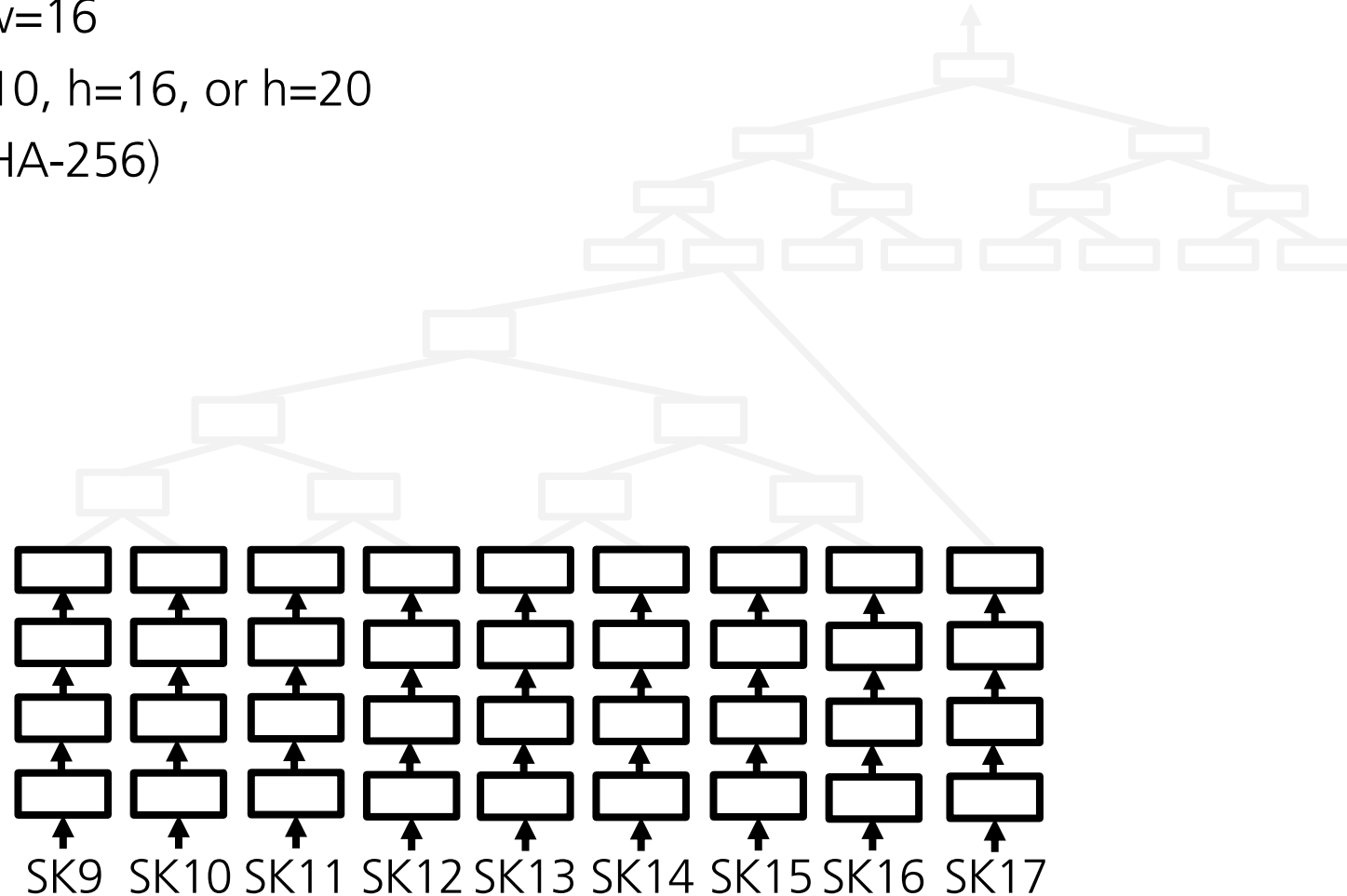
Public Key Generation – WOTS+

IRTF Parameters:

WOTS+ chain length $w=16$

Merkle tree height $h=10$, $h=16$, or $h=20$

256 Bit Hashes (e.g. SHA-256)



Performance Consideration

Public Key Generation – WOTS+

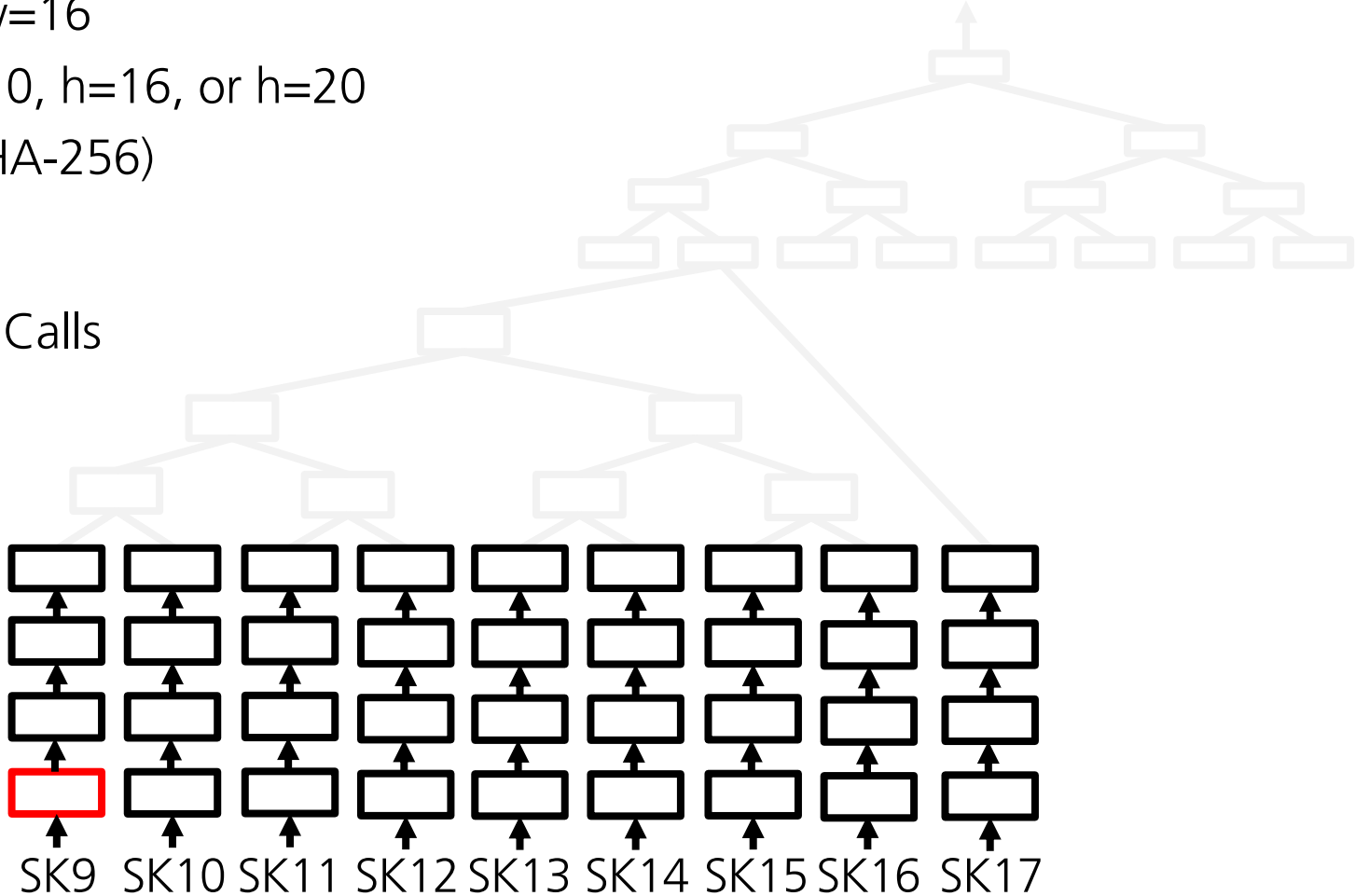
IRTF Parameters:

WOTS+ chain length $w=16$

Merkle tree height $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

 3 Hash Function Calls



Performance Consideration


Public Key Generation – WOTS+

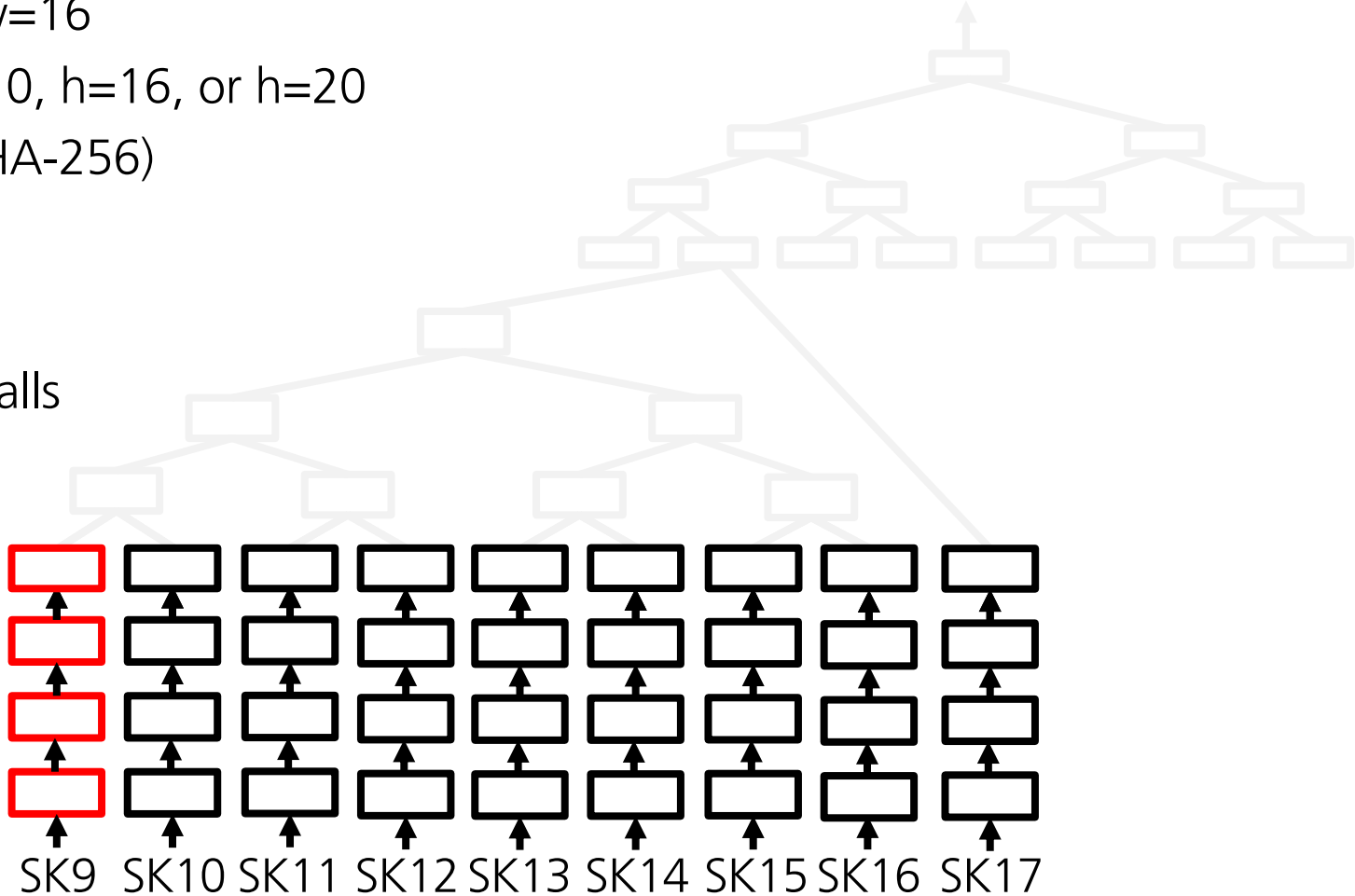
IRTF Parameters:

WOTS+ chain length $w=16$

Merkle tree height $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

 $3*w = 48$
Hash Function Calls



Performance Consideration


Public Key Generation – WOTS+

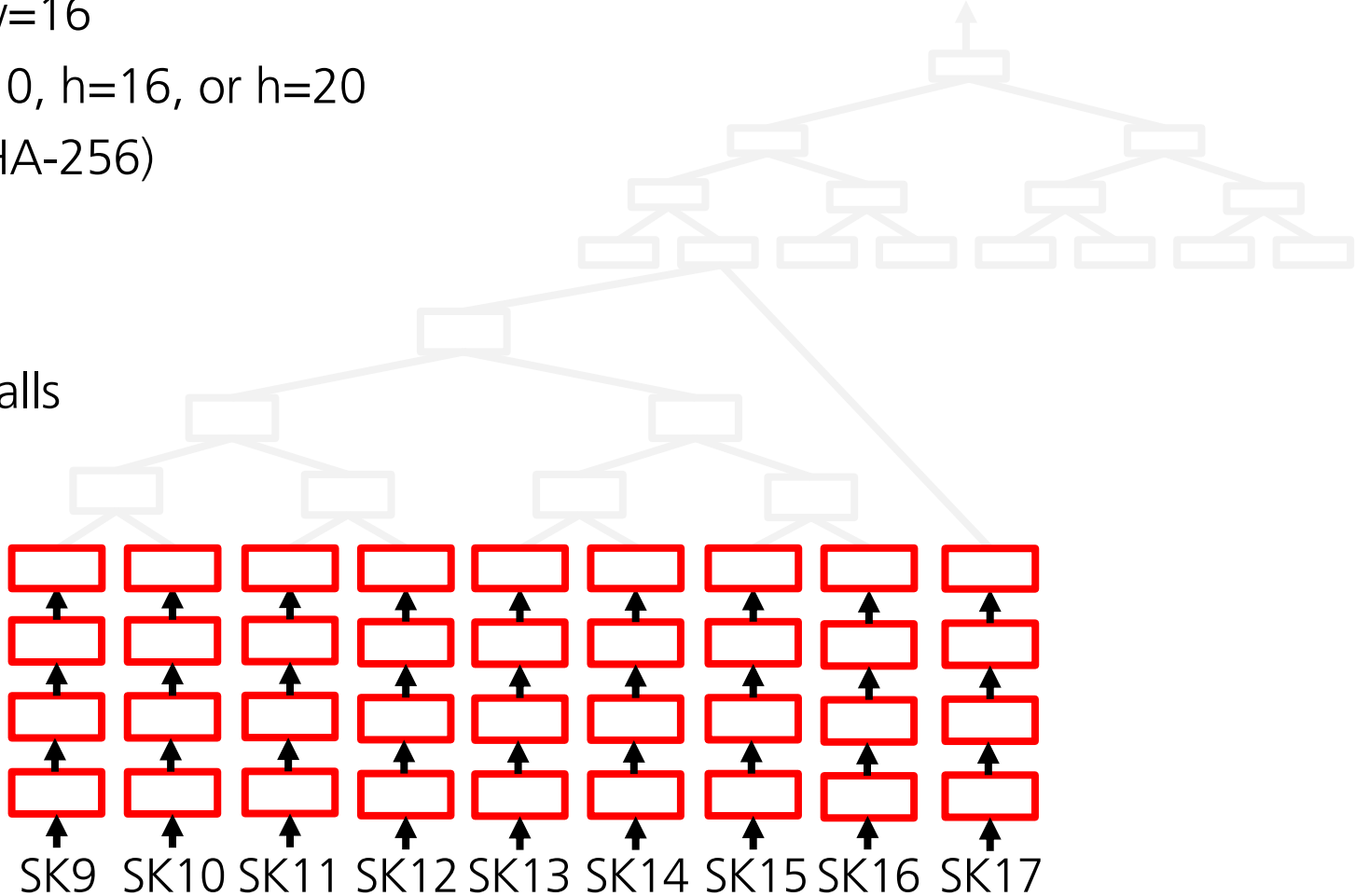
IRTF Parameters:

WOTS+ chain length $w=16$

Merkle tree height $h=10$, $h=16$, or $h=20$

256 Bit Hashes (e.g. SHA-256)

 $48 \cdot 67 = 3216$
Hash Function Calls



Performance Consideration


Public Key Generation – WOTS+

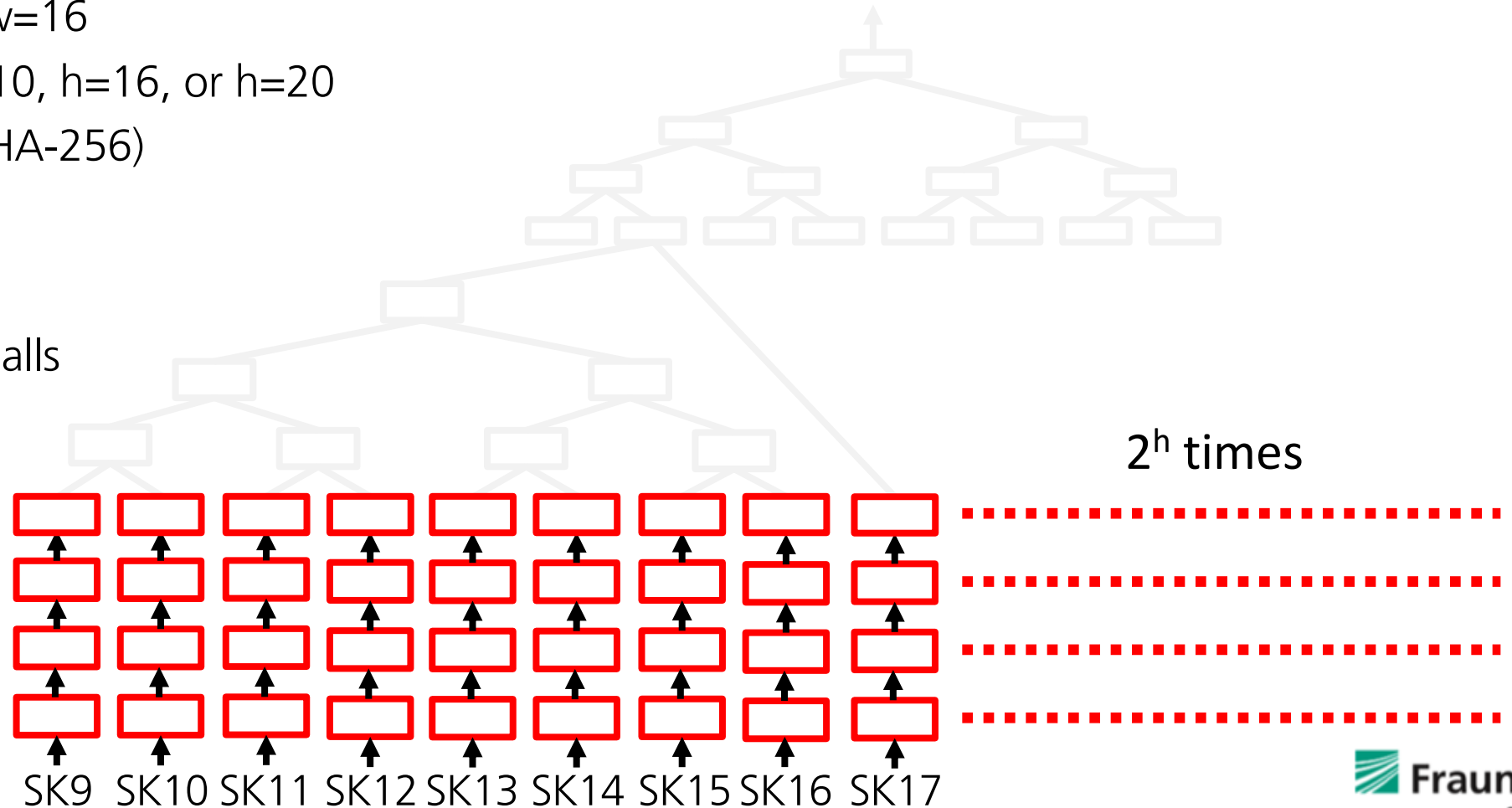
IRTF Parameters:

WOTS+ chain length $w=16$

Merkle tree height $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

 $3216 \cdot 2^h$
Hash Function Calls



Performance Consideration

Public Key Generation – L-Tree

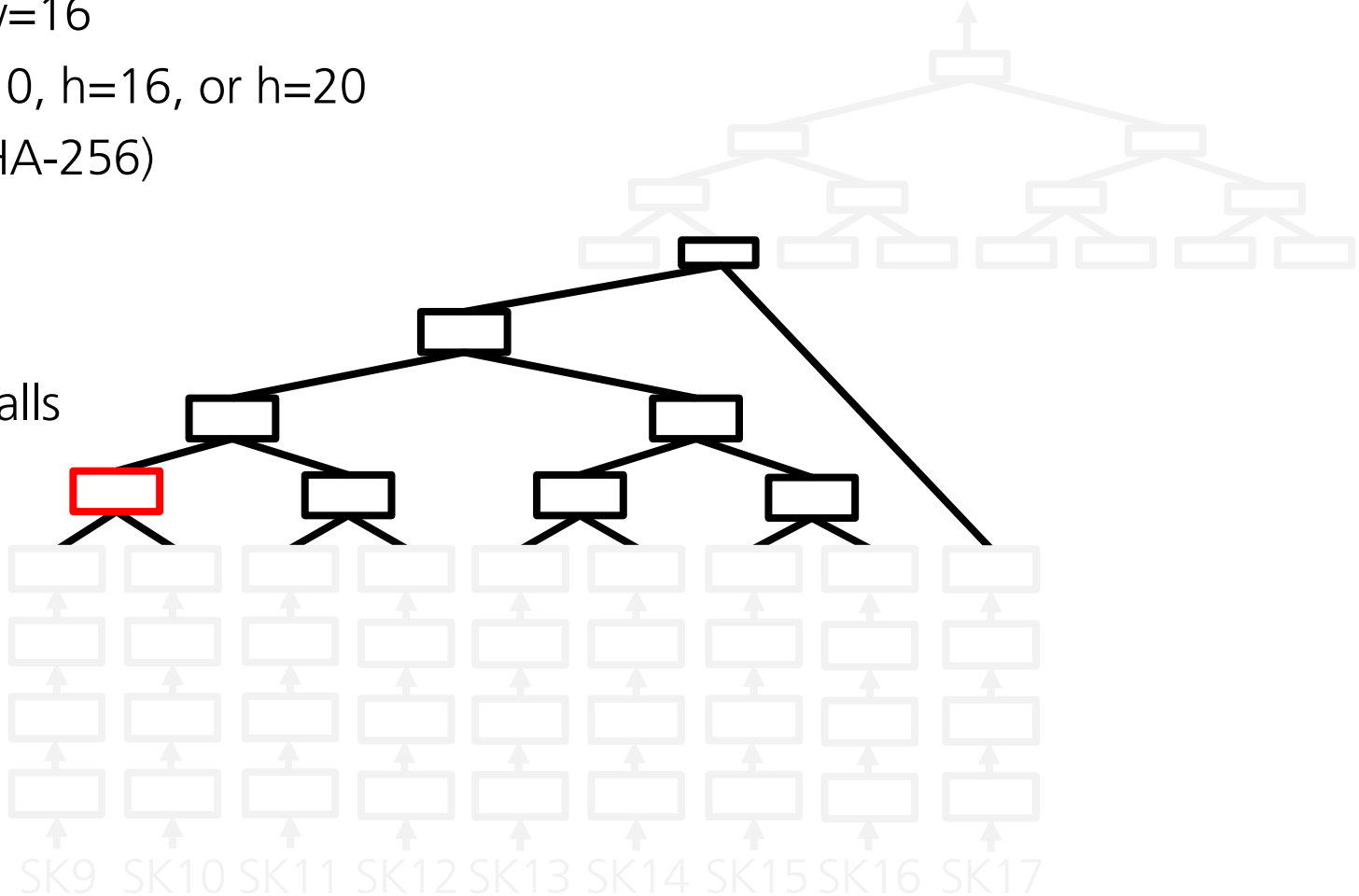
IRTF Parameters:

WOTS+ chain length $w=16$

Merkle tree height $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

 4
Hash Function Calls



Performance Consideration


Public Key Generation – L-Tree

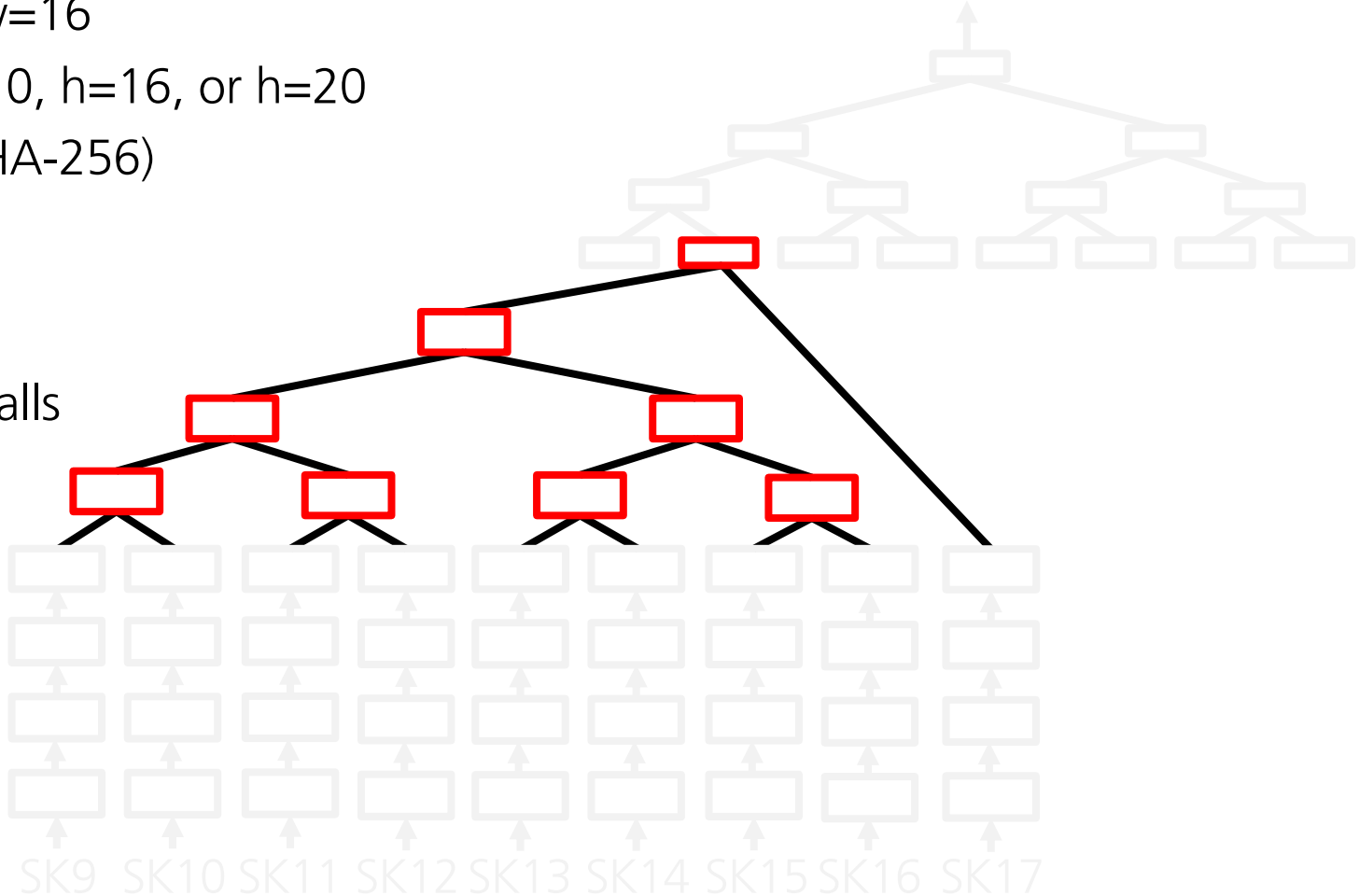
IRTF Parameters:

WOTS+ chain length $w=16$

Merkle tree height $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

 $4 * 65 = 268$
Hash Function Calls



Performance Consideration


Public Key Generation – L-Tree

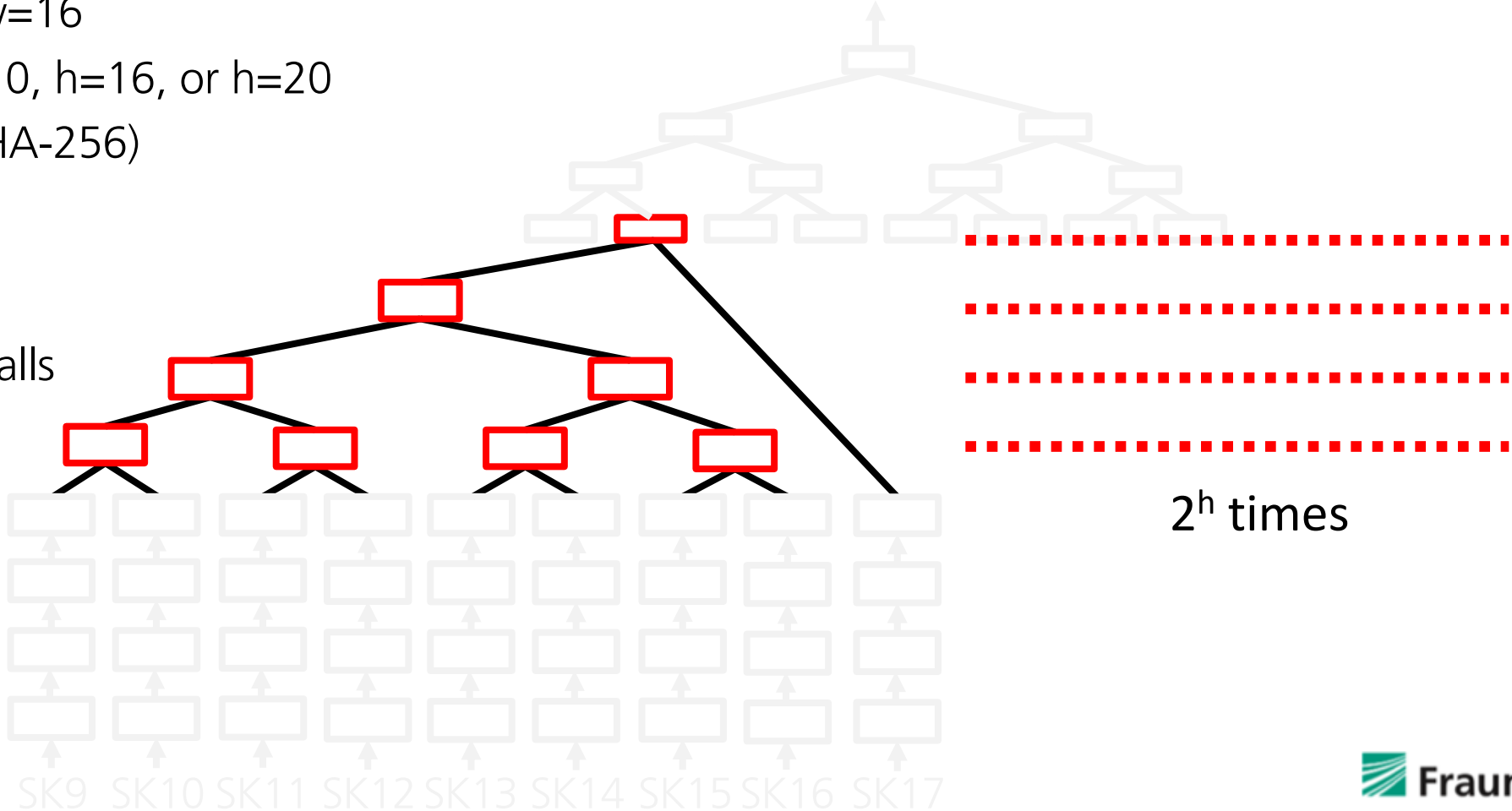
IRTF Parameters:

WOTS+ chain length $w=16$

Merkle tree height $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

 $260 \cdot 2^h$
Hash Function Calls



Performance Consideration


Public Key Generation – XMSS

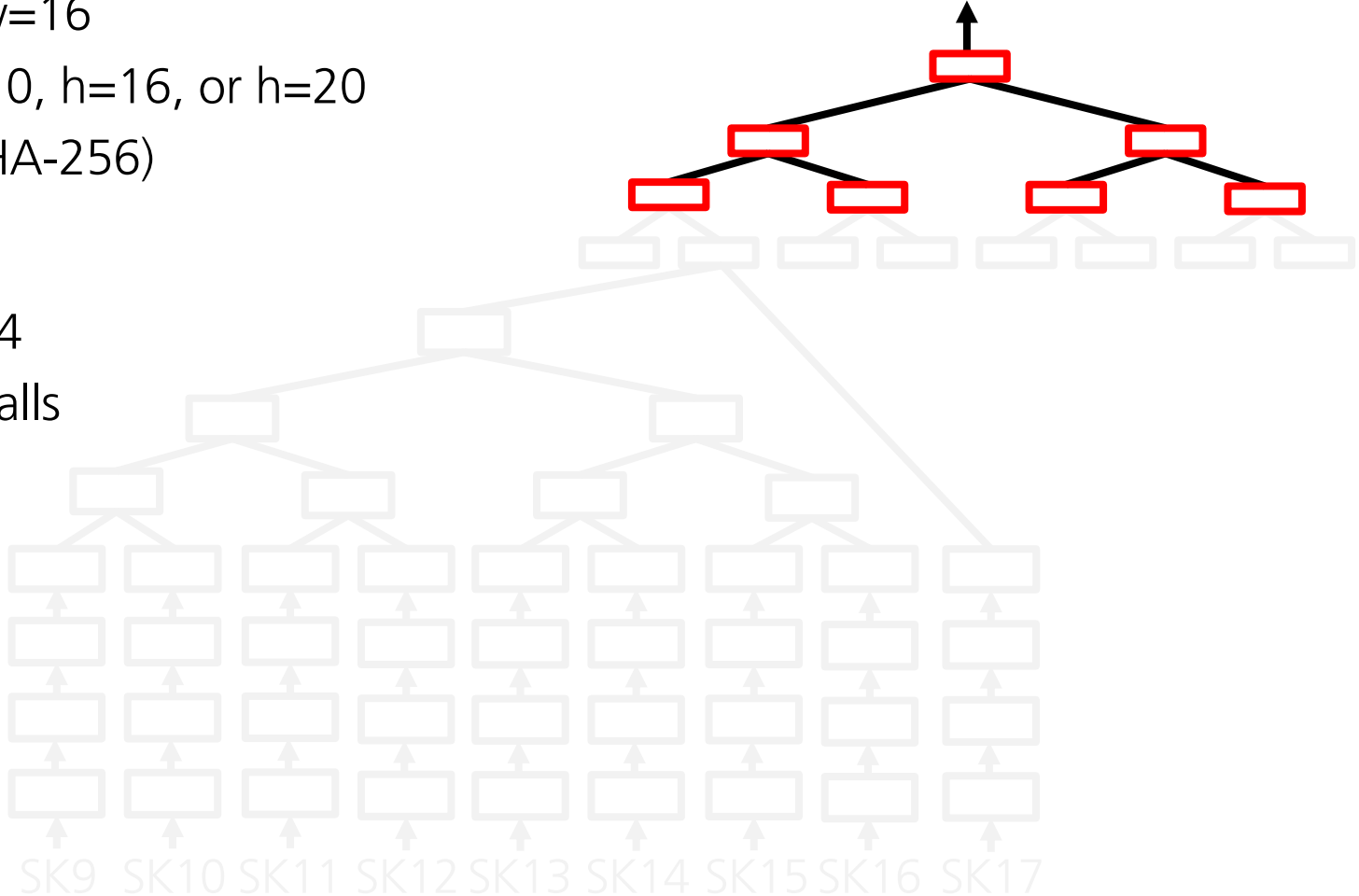
IRTF Parameters:

WOTS+ chain length $w=16$

Merkle tree height $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

 $4 \cdot (2^h - 1) = 4 \cdot 2^h - 4$
Hash Function Calls



Performance Consideration


Public Key Generation – XMSS

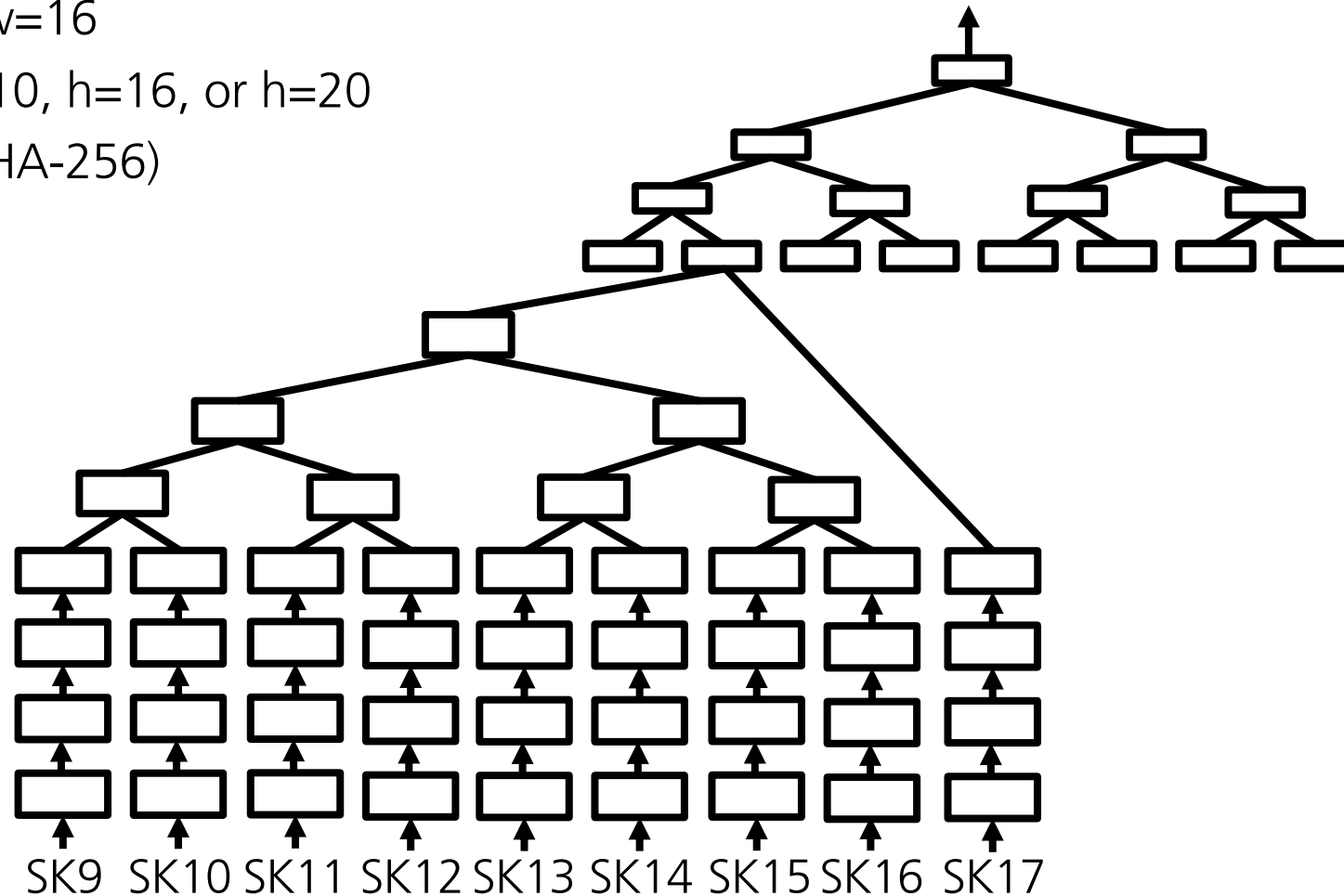
IRTF Parameters:

WOTS+ chain length $w=16$

Merkle tree height $h=10, h=16, \text{ or } h=20$

256 Bit Hashes (e.g. SHA-256)

 $3480 * 2^h - 4$
Total Hash
Function Calls



Performance Consideration

Hash Function Calls

	h=10	h=16	h=20
Signatures	1024	65,536	1,048,576
Public Key Generation	3,563,520	228,065,280	3,649,044,480
Signature Generation	~5,560	~263,684	~4,195,828
Signature Verification	~1,908	~1,932	~1,948

Performance with SHA-256

	h=10	h=16	h=20
Signatures	1024	65,536	1,048,576
Public Key Generation	423,099,648 clock cycles	$27 \cdot 10^9$ clock cycles	$434 \cdot 10^9$ clock cycles
With 400 MHz	<1.1 s	<70 s	<1085 s
Sign	< 2 ms	< 70 ms	< 1 s
Verify	< 1 ms	< 1 ms	< 1 ms

Performance with SHA-3

	h=10	h=16	h=20
Signatures	1024	65,536	1,048,576
Public Key Generation	79,159,200 clock cycles	$5 * 10^9$ clock cycles	$81 * 10^9$ clock cycles
With 400 MHz	< 200 ms	< 12.5 s	< 203 s
Sign	< 1 ms	< 12.5 ms	< 200 ms
Verify	< 1 ms	< 1 ms	< 1 ms

Comparison with ECC

FPGA Implementation Estimates (Virtex-5)

	Ed25519	XMSS-SHA3 h=10
Public Key Generation	< 1 ms	< 200 ms
Sign	< 1 ms	< 1 ms
Verify	< 2 ms	< 1 ms

Optimisations and Trade-Offs

Parallelization and Caching

- Parallelization
 - WOTS+ trivial to compute in parallel
 - L-Tree and XMSS more difficult to parallelize
- More/Less Caching
 - More caching of XMSS for authentication path (costs more memory)
 - ➔ Improves the signing performance
 - Less caching to save memory
 - ➔ In the worst case, signing almost as slow as public key generation
 - ➔ Useful for lightweight applications with low memory



Thank you for your attention!